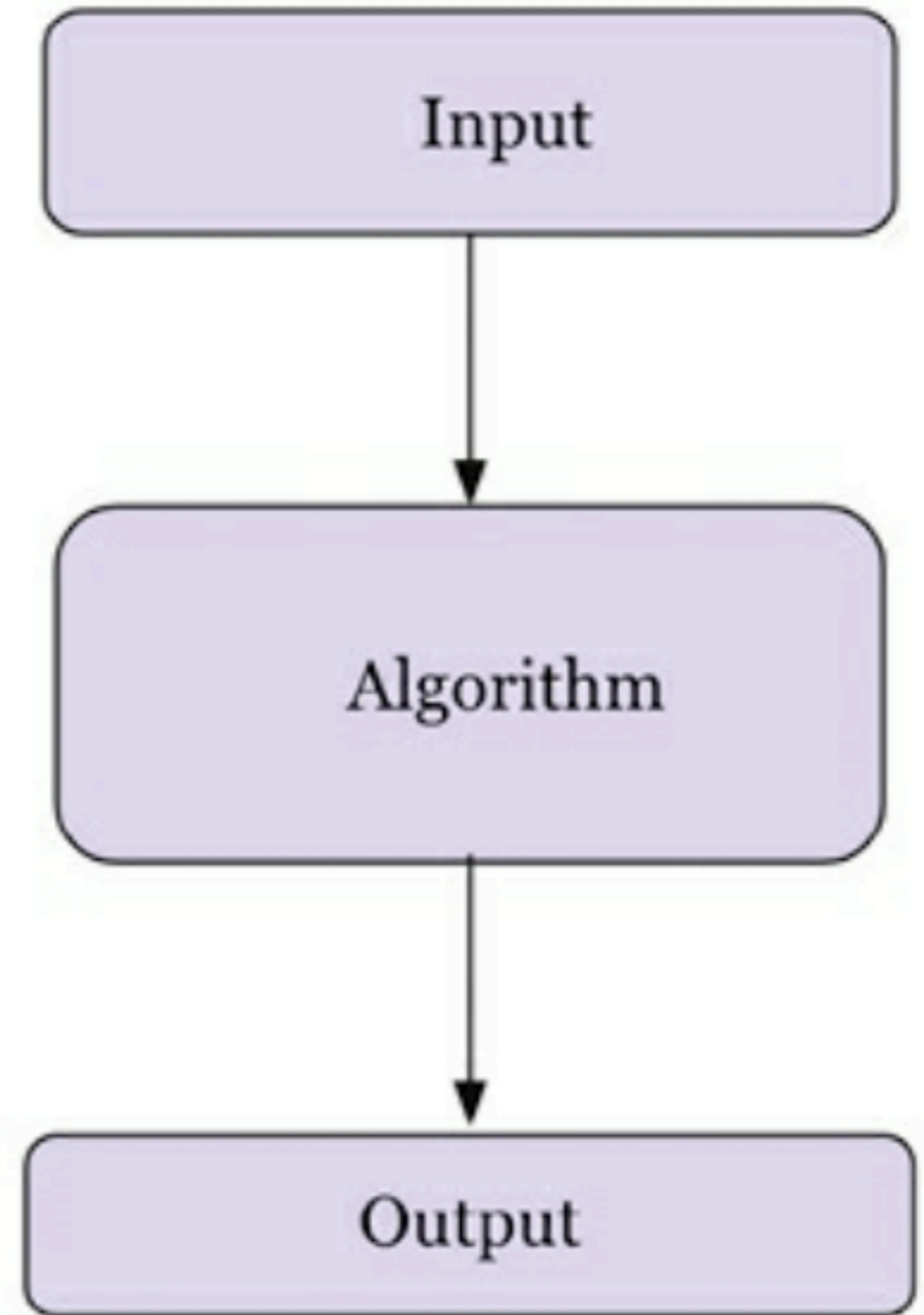


2 Introduction to Algorithm Design

For COMSC 132

Algorithms

- Well-defined procedure
- Takes input data
- Process it
- Produces desired output



Performance Analysis

- Time complexity
 - Time required to perform task
- Space complexity
 - Memory required to perform task

Linear in n

- As size of problem increases
- Time required grows linearly

```
import time

for power in range(3, 9):
    n = 10 ** power
    t0 = time.time()
    total = 0
    for i in range(n):
        total += 1
    elapsed = round(time.time() - t0, 4)
    print("{0:>5}".format(elapsed), n)
```

```
⇒ 0.0002 1000
   0.0018 10000
   0.0141 100000
   0.1177 1000000
   1.0664 10000000
  11.5965 100000000
```

```
import time
```

```
for power in range(3, 9):
    n = 10 ** power
    t0 = time.time()
    total = 0
    for i in range(n):
        total += 1
    elapsed = round(time.time() - t0, 4)
    print("{0:>5}".format(elapsed), n)
```

Asymptotic notation

- Θ (Theta)
 - Worst-case running time with a tight bound
- O (Big Oh)
 - Worst-case with an upper bound
- Ω (Omega)
 - Lower bound of running time

- Quadratic
 $O(n^2)$
- Cubic
 $O(n^3)$
- Exponential
 $O(2^n)$

Time Complexity	Name
$O(1)$	Constant
$O(\log n)$	Logarithmic
$O(n)$	Linear
$O(n \log n)$	Linear-logarithmic
$O(n^2)$	Quadratic
$O(n^3)$	Cubic
$O(2^n)$	Exponential

Table 2.1: Runtime complexity of different functions

Take Largest Term

- We only care about case for large n
- So if
 - Time = $4n^{**3} + 100n + 1000$
 - Complexity is $O(n^{**3})$

$O(n^{**2})$ Complexity

```
✓ [11] import time
12s
    for power in range(2, 5):
        n = 10 ** power
        t0 = time.time()
        total = 0
        for i in range(n):
            for j in range(n):
                total += 1
        elapsed = round(time.time() - t0, 4)
        print(f'{elapsed:9.4f}', f'{n:>7,}')
```

```
⇒ 0.0015    100
   0.1825   1,000
  12.2375  10,000
```

```
import time
```

```
for power in range(2, 5):
    n = 10 ** power
    t0 = time.time()
    total = 0
    for i in range(n):
        for j in range(n):
            total += 1
    elapsed = round(time.time() - t0, 4)
    print(f'{elapsed:9.4f}', f'{n:>7,}')
```


$O(2^{**n})$ Complexity

```
import time

for n in range(15, 25):
    t0 = time.time()
    total = 0
    for i in range(2**n):
        total += 1
    elapsed = round(time.time() - t0, 4)
    print(f'{elapsed:9.4f}', f'{n:>7,}')
```

↩	0.0048	15
	0.0067	16
	0.0187	17
	0.0277	18
	0.0557	19
	0.1086	20
	0.2344	21
	0.4506	22
	0.8724	23
	1.7665	24

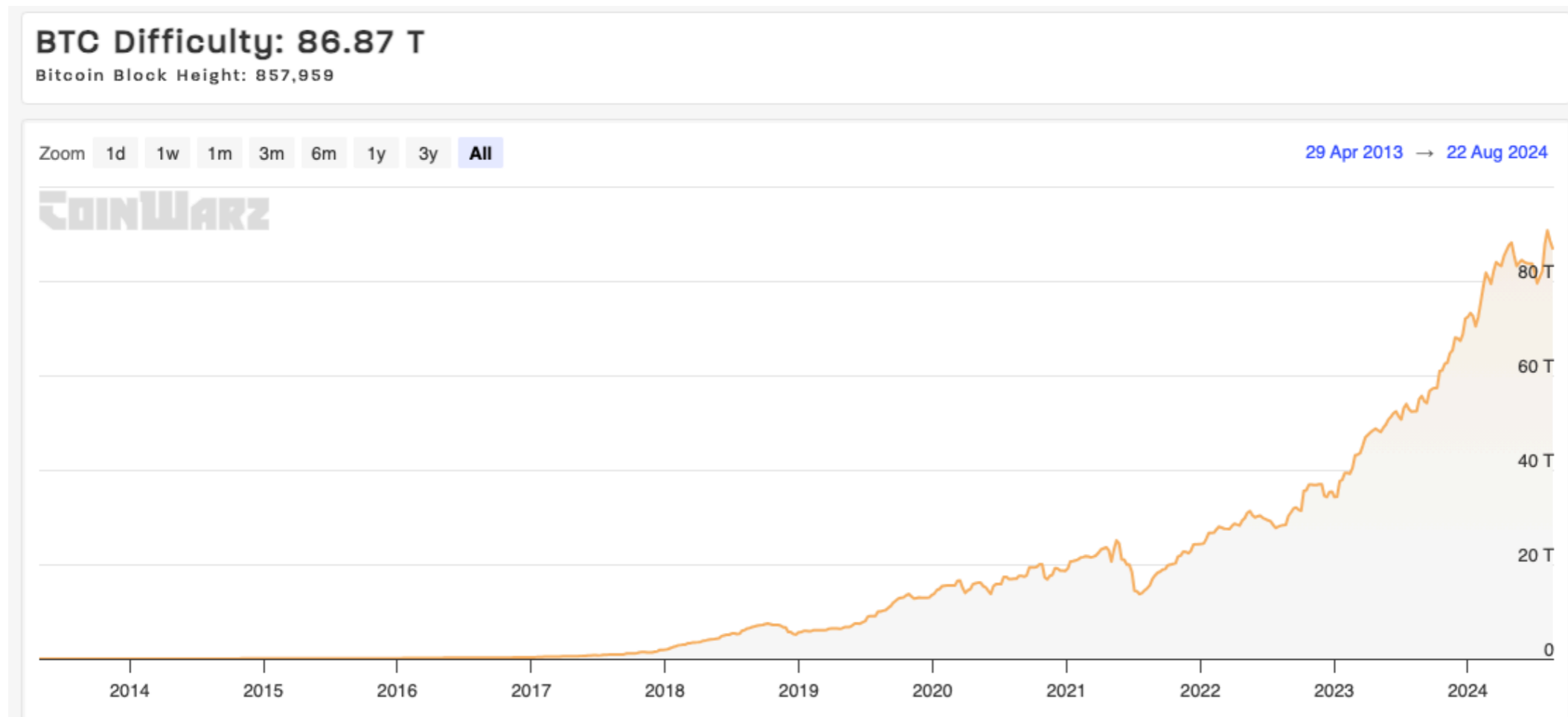
```
import time

for n in range(15, 25):
    t0 = time.time()
    total = 0
    for i in range(2**n):
        total += 1
    elapsed = round(time.time() - t0, 4)
    print(f'{elapsed:9.4f}', f'{n:>7,}')
```

Bitcoin

- Miners take a block of transactions
 - Add a random "nonce"
 - Calculate $\text{SHA256}(\text{SHA256}(\text{block}))$
 - If (hash < **target**):
 - Win! You get a reward.
 - else:
 - Go pick a new nonce

Bitcoin Difficulty

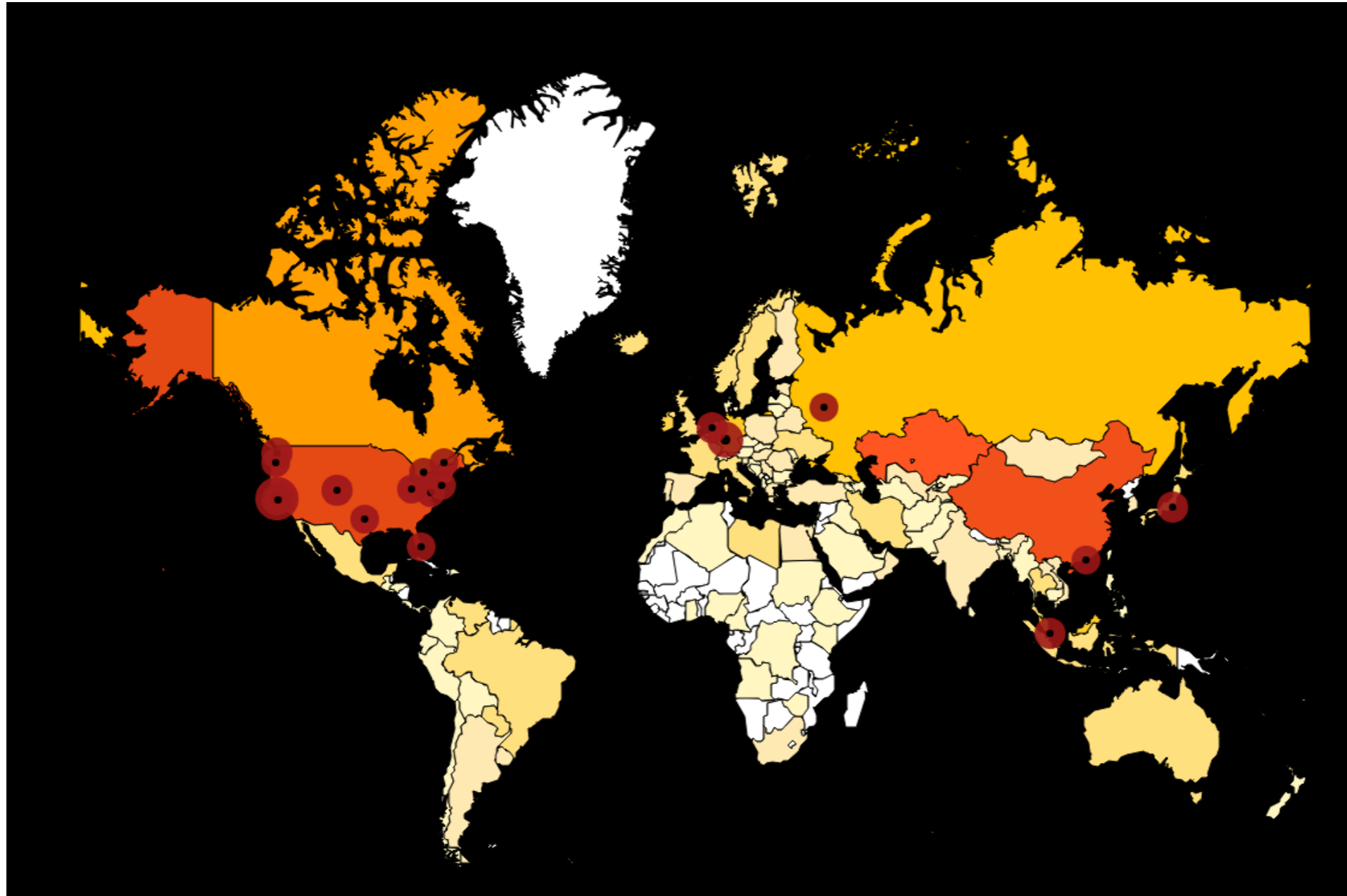


- From <https://www.coinwarz.com/mining/bitcoin/difficulty-chart>

Bitcoin is CPU-Hard

- Requires lots of CPU cycles to calculate SHA256
- Ultimate cost is power
- Miners move to locations with cheap electricity

What Countries Mine Bitcoin?



- <https://chainbulletin.com/bitcoin-mining-map/>

Country overall hash power in EH/s



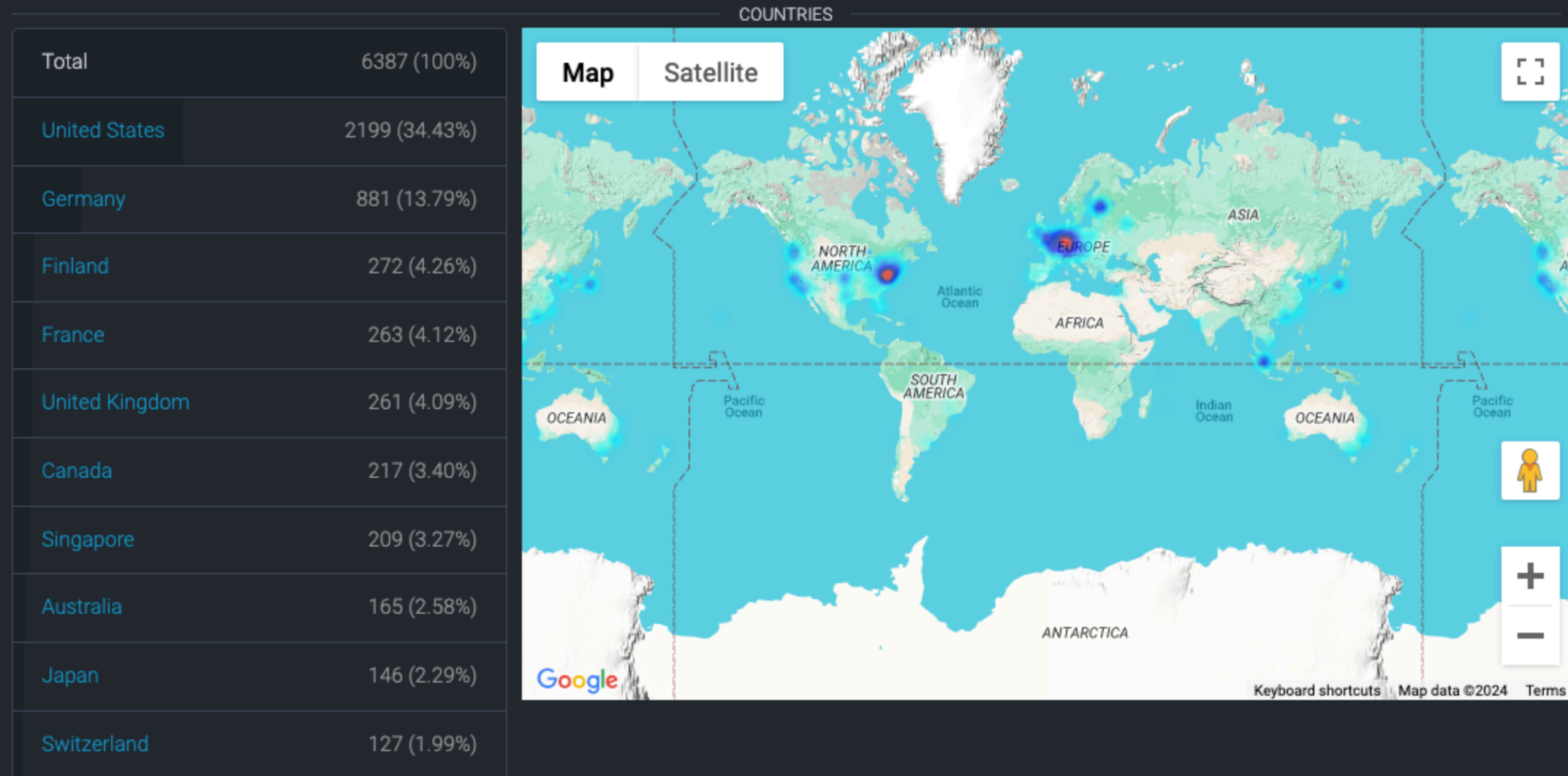
Country hashrate as percentage from total output

Ethereum is Memory-Hard

- Mining requires a data structure called a **Directed Acyclic Graph (DAG)**
- DAG must be generated in each epoch
- The DAG is currently 5 GB
- It grows with each epoch
 - <https://depoli.ca/ethereum-algorithm>

Ethereum Mainnet Statistics

Clients Countries Sync Status OS Network Types History



- <https://www.ethernodes.org/countries>

Kahoot!

Ch 2