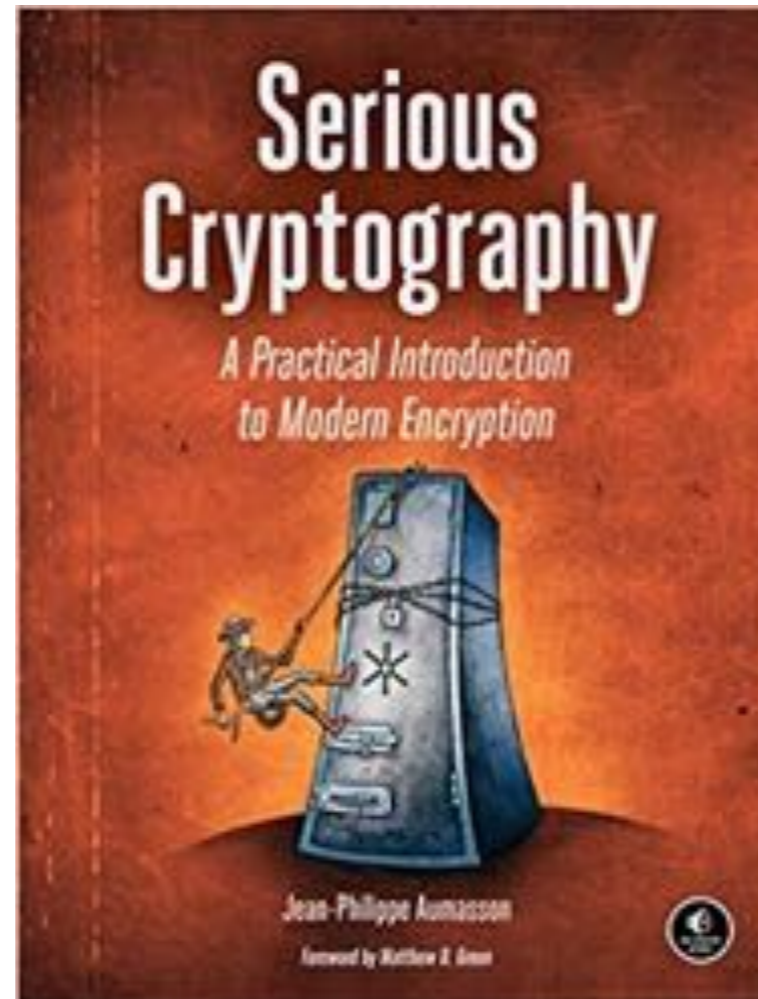


# CNIT 141

## Cryptography for Computer Networks



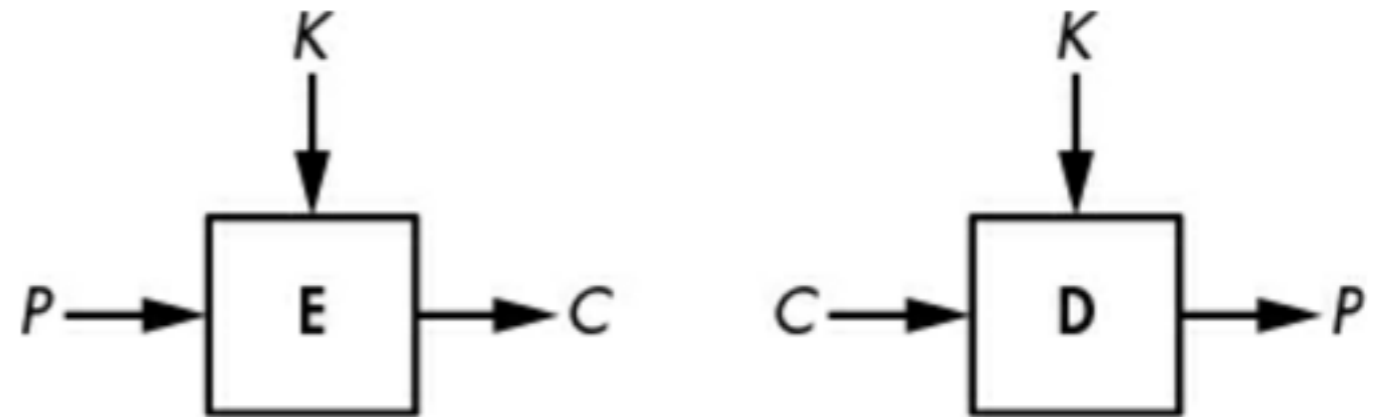
### 1. Encryption

# Topics

- The Basics
- Classical Ciphers
- How Ciphers Work
- Perfect Encryption: The One-Time Pad
- Encryption Security
- Asymmetric Encryption
- When Ciphers Do More Than Encryption
- How Things Can Go Wrong

# The Basics

- P: Plaintext
- K: Key
- C: Cleartext
- E: Encryption via cipher
- D: Decryption via cipher



*Figure 1-1: Basic encryption and decryption*

# Classical Ciphers

# Caesar Cipher

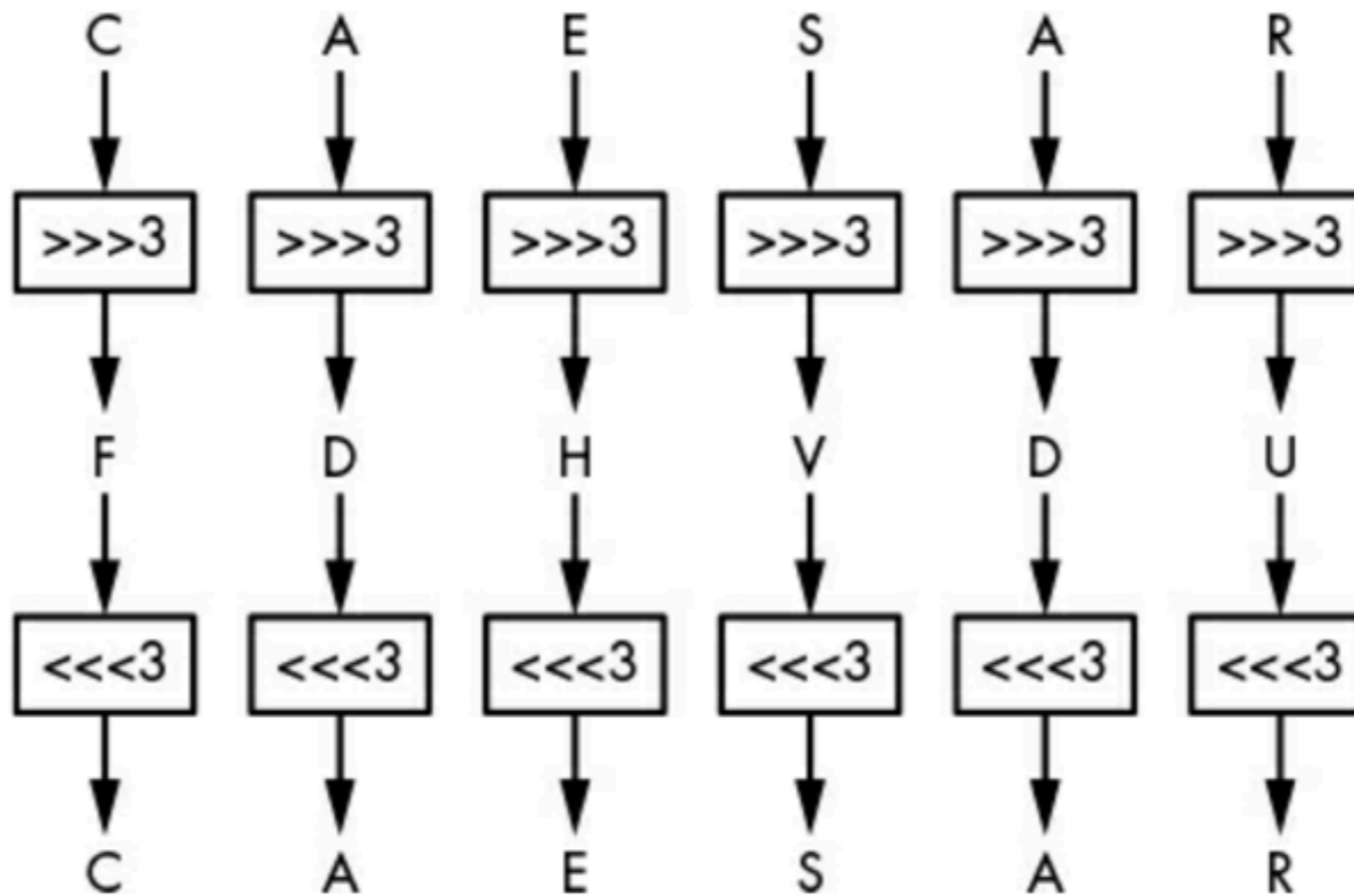


Figure 1-2: The Caesar cipher

# Caesar Cipher in Python

```
GNU nano 2.0.6                               File: caesar1.py
letters = "ABCDEFGHIJKLMNOPQRSTUVWXYZ"
P = raw_input("Plaintext: ")
shift = int(raw_input("Shift: "))
C = ""
for c in P:
    i = letters.find(c) + shift
    if i >= len(letters):
        i = i - len(letters)
    C = C + letters[i]
print "Ciphertext: ", C
```

```
Sams-MacBook-Pro:lec sambowne$ python caesar1.py
Plaintext: HELLO
Shift: 3
Ciphertext:  KHOOR
```

# Brute Force Attack in Python

```
GNU nano 2.0.6                               File: caesar2.py

ciphertext = raw_input("Ciphertext: ")

letters = "ABCDEFGHIJKLMNOPQRSTUVWXYZ"

for shift in range(26):
    cleartext = ""
    for c in ciphertext:
        i = letters.find(c) + shift
        if i >= len(letters):
            i = i - len(letters)
        cleartext = cleartext + letters[i]
    print shift, cleartext
```

# Brute Force Attack in Python

```
Sams-MacBook-Pro:lec sambowne$ python caesar2.py  
Ciphertext: KHOOR  
0 KHOOR  
1 LIPPS  
2 MJQQT  
3 NKRRU  
4 OLSSV  
5 PMTTW  
6 QNUUX  
7 ROVVY  
8 SPWWZ  
9 TQXXA  
10 URYYB  
11 VSZZC  
12 WTAAD  
13 XUBBE  
14 YVCCF  
15 ZWDDG  
16 AXEEH  
17 BYFFI  
18 CZGGJ  
19 DAHHK  
20 EB IIL  
21 FCJJM  
22 GDKKN  
23 HELLO  
24 IFMMP  
25 JGNNQ
```



# Vigenere Cipher

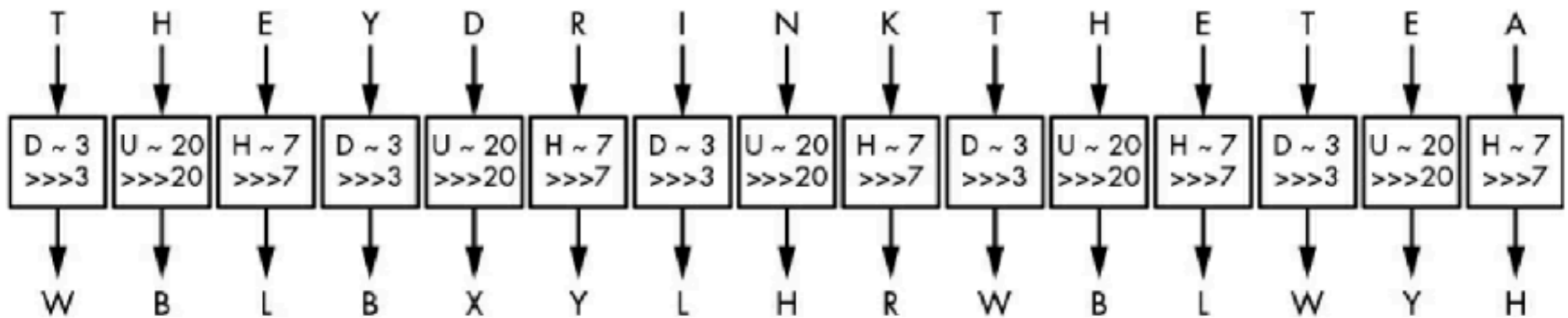


Figure 1-3: The Vigenère cipher

- Shift varies with a repeated keyword
- Combine several Caesar ciphers together

# Breaking the Vigenere Cipher

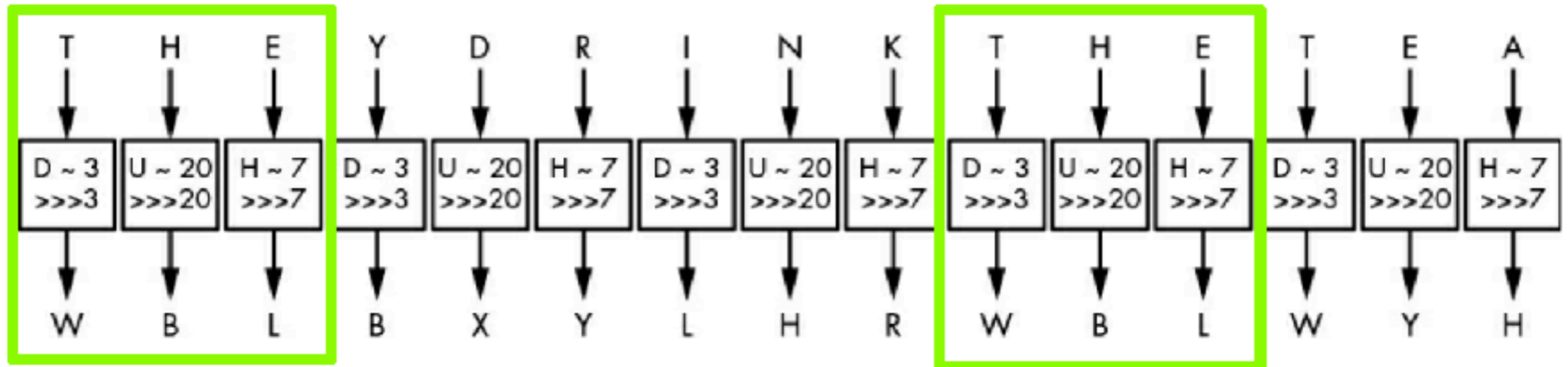
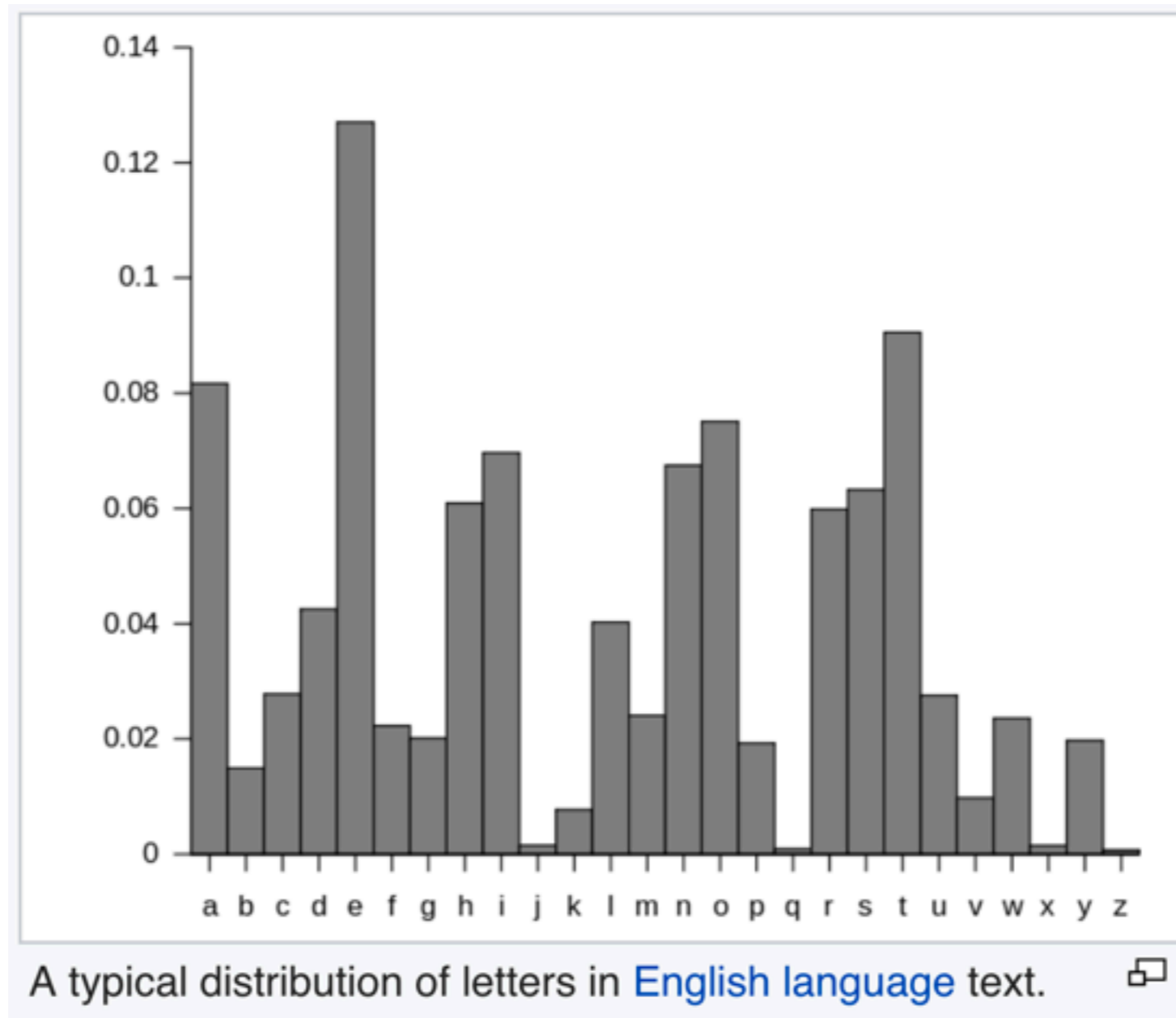


Figure 1-3: The Vigenère cipher

- Find repeating ciphertext to deduce key length
- Use *frequency analysis*

# Frequency Analysis



- From Wikipedia

# Modified Caesar Program

- Converts to uppercase
- Preserves spaces

```
GNU nano 2.0.6                               File: caesar1a.py
letters = "ABCDEFGHIJKLMNOPQRSTUVWXYZ"
P = raw_input("Plaintext: ").upper()
shift = int(raw_input("Shift: "))
C = ""
for c in P:
    if c != " ":
        i = letters.find(c) + shift
        if i >= len(letters):
            i = i - len(letters)
        C = C + letters[i]
    else:
        C = C + " "
print "Ciphertext: ", C
```

# Encrypt a Paragraph

```
Sams-MacBook-Pro:lec sambowne$ python caesar1a.py  
Plaintext: Four score and seven years ago our fathers brought forth on this continent, a new nation, conceived in Liberty, and dedicated to the proposition that all men are created equal.  
Shift: 3  
Ciphertext: IRXU VFRUH DQG VHYHQ BHDUV DJR RXU IDWKHUV EURXJKW IRUWK RQ WKLW FR  
QWLQHQC D QHZ QDWLRQC FRQFHLYHG LQ OLEHUWBC DQG GHGLFDWHG WR WKH SURSRVLWLRQ WK  
DW DOO PHQ DUH FUHDWHG HTXDOC  
Sams-MacBook-Pro:lec sambowne$
```

# Frequency Counter

```
GNU nano 2.0.6                               File: freq1.py
letters = "ABCDEFGHIJKLMNOPQRSTUVWXYZ"
P = raw_input("Text: ").upper()

count = [0]*26

for c in P:
    if c != " ":
        i = letters.find(c)
        count[i] += 1

for c in letters:
    i = letters.find(c)
    print c, count[i]
```

```
Sams-MacBook-Pro:lec sambowne$ python freq1.py
Text: IRXU VFRUH DQG VHYHQ BHDUV DJR RXU IDWKHUV EURXJKW IRUWK RQ WKLW FRQWLQHQW
C D QHZ QDWLRQC FRQFHLYHG LQ OLEHUWBC DQG GHGLFDWHG WR WKH SURSRVLWLRQ WKDW DOO
PHQ DUH FUHDWHG HTXDOC
A 0
B 2
C 4
D 13
E 2
F 6
G 7
H 18
I 3
J 2
K 6
L 9
M 0
N 0
O 4
P 1
Q 14
R 14
S 2
T 1
U 11
V 6
W 15
X 4
Y 2
Z 1
```

# How Ciphers Work



# Two Components

- **Permutation**

- Transforms one letter to another letter
- In Caesar cipher, shift letter three places

- **Mode of Operation**

- Algorithm to handle messages of arbitrary size
- In Caesar cipher, process each letter independently

# Permutation Security

- **Permutation should be determined by the key**
  - If key is secret, attacker can't easily decrypt
- **Different keys should result in different permutations**
- **Permutation should look random**
  - No pattern in ciphertext

# Mode of Operation

- Caesar cipher encrypts letters one at a time
  - Double letters remain doubled
  - HELLO -> KHOOR
- Patterns in plaintext are preserved in ciphertext
- Insecure (now called “Electronic Code Book” mode)
- More secure modes encrypt repeated text differently each time

# Perfect Encryption: The One-Time Pad

# XOR

- XOR combines two bits
  - $0 \wedge 0 = 0$
  - $0 \wedge 1 = 1$
  - $1 \wedge 0 = 1$
  - $1 \wedge 1 = 0$

# Encrypting a Stream of Bits

- Plain: ABC = 0100 0001 0100 0010 0100 0011
  - Key: 0110 0110 0110 0101 1010 1110
  - Cipher: 0010 0111 0010 0111 1110 1101
- 
- Key must be random and never re-used
  - Key must be longer than all the plaintexts you want to send

# Unbreakable

- If an attacker uses a brute-force attack
- Trying all possible keys
- They get all possible letter sequences
- No way to identify the correct decryption

**Kahoot!**

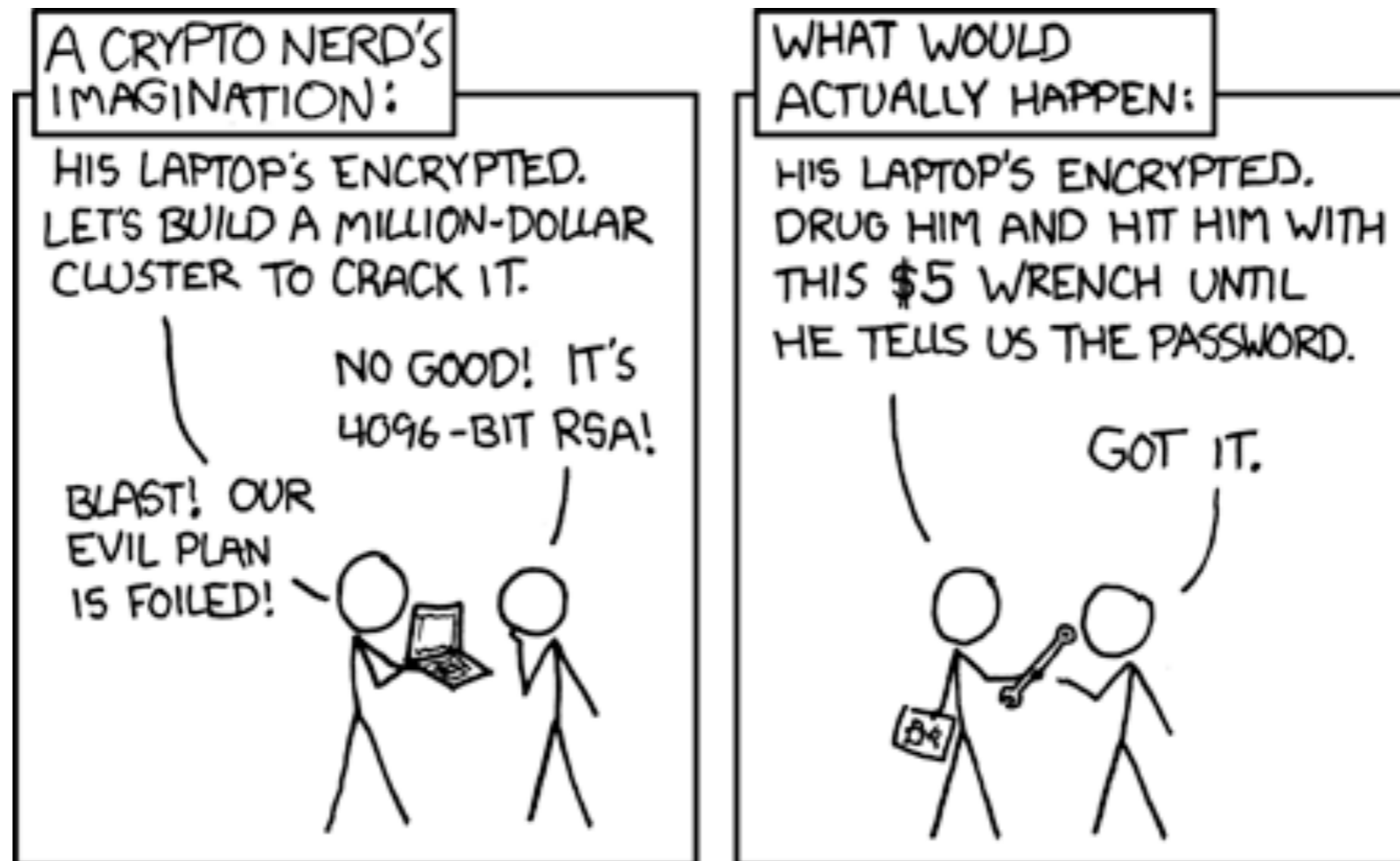


# Encryption Security

# Attack Models

- Set requirements for cryptographers who design ciphers
  - So they know what attacks to prevent
- Give guidelines to users
  - Whether a cipher is safe in their environment
- Provide clues for cryptanalysts who attempt to break ciphers
  - Is an attack doable in the model considered?

# Attack Models



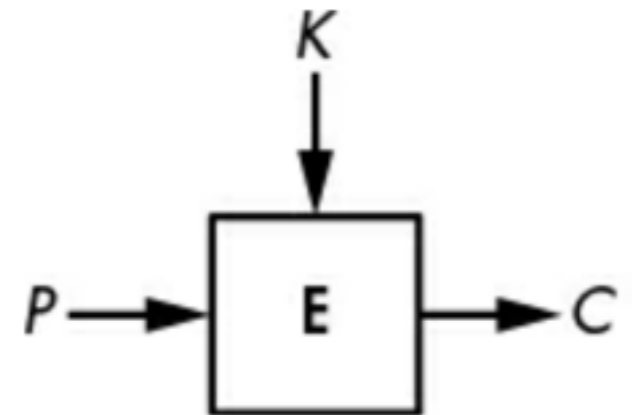
# Kerckhoff's Principle

- The **key** is secret
- The **cipher** is not secret

# Black-Box Models

## No knowledge of cipher operation

- Ciphertext-Only Attack (COA)
  - Attacker sees only  $C$
- Known-Plaintext Attack (KPA)
  - Attacker knows  $P$  and  $C$
- Chosen-Plaintext Attack (CPA)
  - Attacker can perform encryption for any  $P$
- Chosen-Ciphertext Attack (CCA)
  - Attacker can perform encryption and decryption



# Gray-Box Models

- Attacker has access to the implementation
  - Can tamper with the system's internals
- **Side-channel attack**
  - Attacker measures something else about the cipher's operation
  - Such as timing or power consumption
  - **Noninvasive** — does not alter integrity of system

# Gray-Box Models

- **Invasive attacks**
  - Modify system
  - Examples
    - Using acid to dissolve parts of a microchip
    - Injecting faults with lasers

# Security Goals

- **Indistinguishability**
  - Ciphertext should be indistinguishable from a random string
- **Non-malleability**
  - Ciphertext cannot be altered and produce meaningful plaintext



# Security Notions

- **IND-CPA**
  - Indistinguishability against a Chosen-Plaintext Attack
  - Also called **semantic security**
  - Two identical plaintext strings must result in different ciphertexts
  - Accomplished by adding “random” bits each time you encrypt

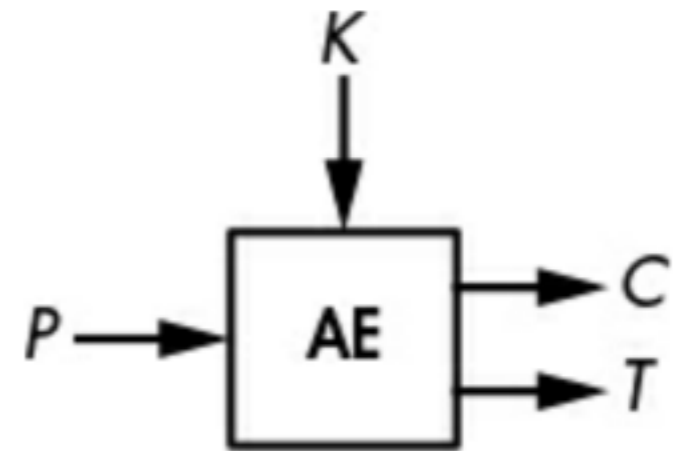
# Asymmetric Encryption

- Uses two keys
- Also called Public-Key encryption
- **Public key** freely published to everyone
- **Private key** held secret
- Will be in later chapters—everything in this chapter is about **symmetric encryption**

# When Ciphers Do More Than Encryption

# Authenticated Encryption

- Returns an **authentication tag** with the ciphertext
- Tag ensures integrity of the message and also authenticates the author
- **Authenticated Encryption with Associated Data (AWAD)**
  - Another variant



*Figure 1-4: Authenticated encryption*

# Format-Preserving Encryption

- Normally encryption takes inputs as bits and returns outputs as bits
- Could be written as hex, base64, etc.
- Format-Preserving Encryption returns ciphertext in the same format as the plaintext
  - Zip code -> Zip code
  - IP address -> IP address

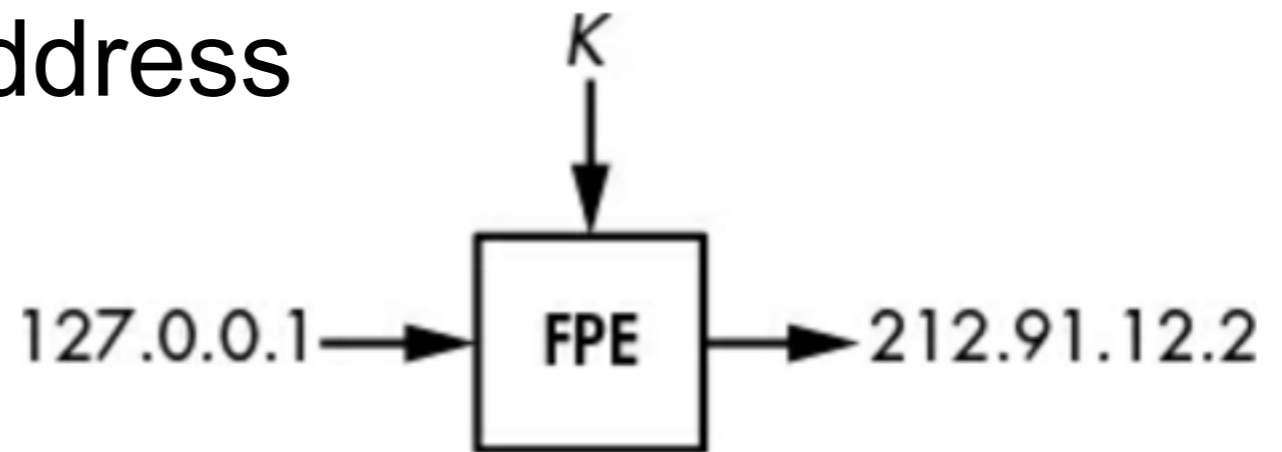


Figure 1-5: Format-preserving encryption for IP addresses

# Fully Homomorphic Encryption

- Allows modification of encrypted data without decrypting it
- The first FHE scheme was created in 2009
- Very slow

# Searchable Encryption

- Searches encrypted data without decrypting it
- Using an encrypted search string
- Protects privacy of search engine users
- Experimental at present

# Tweakable Encryption

- Adds a “tweak” parameter to normal encryption
  - Such as a unique customer number
- Main application is disk encryption

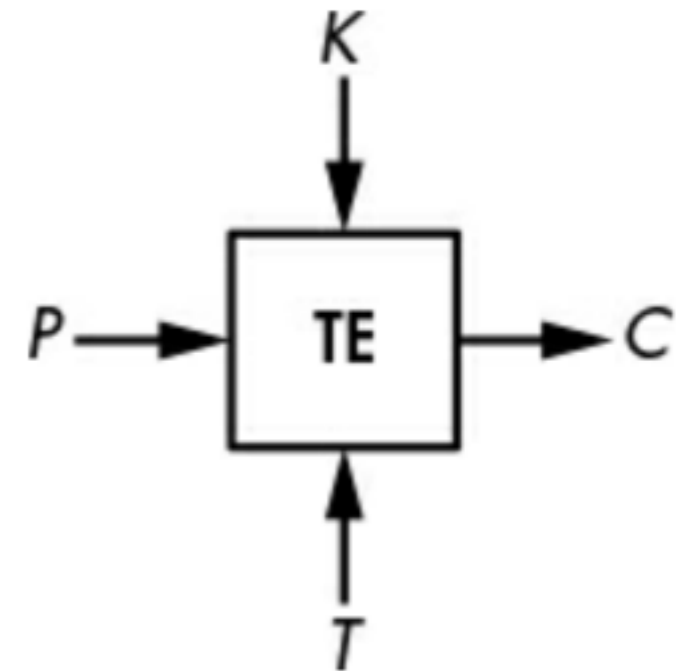


Figure 1-6: Tweakable encryption



# How Things Can Go Wrong

# Weak Cipher

- 2G phone networks used the A5/1 cipher
- Vulnerable to a **time-memory trade-off** attack
  - Using large lookup tables to speed up an attack

# Wrong Model

- Padding Oracle attack
- If a user submitted data that decrypted to a valid string, that was taken as authentication
  - Even if the string contained nonsense
- Server provided error messages for incorrect padding
- Those errors can be used to find valid ciphertext without knowing the key

**Kahoot!**