Windows Internals

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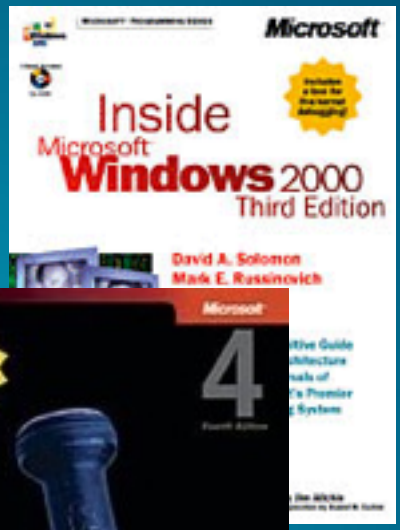
About the Speaker: David Solomon

- 1982-1992: VMS operating systems development at Digital
- 1992-present: Researching, writing, and teaching Windows operating system internals
- Frequent speaker at technical conferences (Microsoft TechEd, IT Forum, PDCs, …)
- Microsoft Most Valuable Professional (1993, 2005)
- Books
  - *Windows Internals, 4th edition*
    - PDF version ships with Server 2003 Resource Kit
  - *Inside Windows 2000, 3rd edition*
  - *Inside Windows NT, 2nd edition*
  - *Windows NT for OpenVMS Professionals*
- Live Classes
  - 2-5 day classes on Windows Internals, Advanced Troubleshooting
- Video Training
  - 12 hour interactive internals tutorial
  - Licensed by MS for internal use



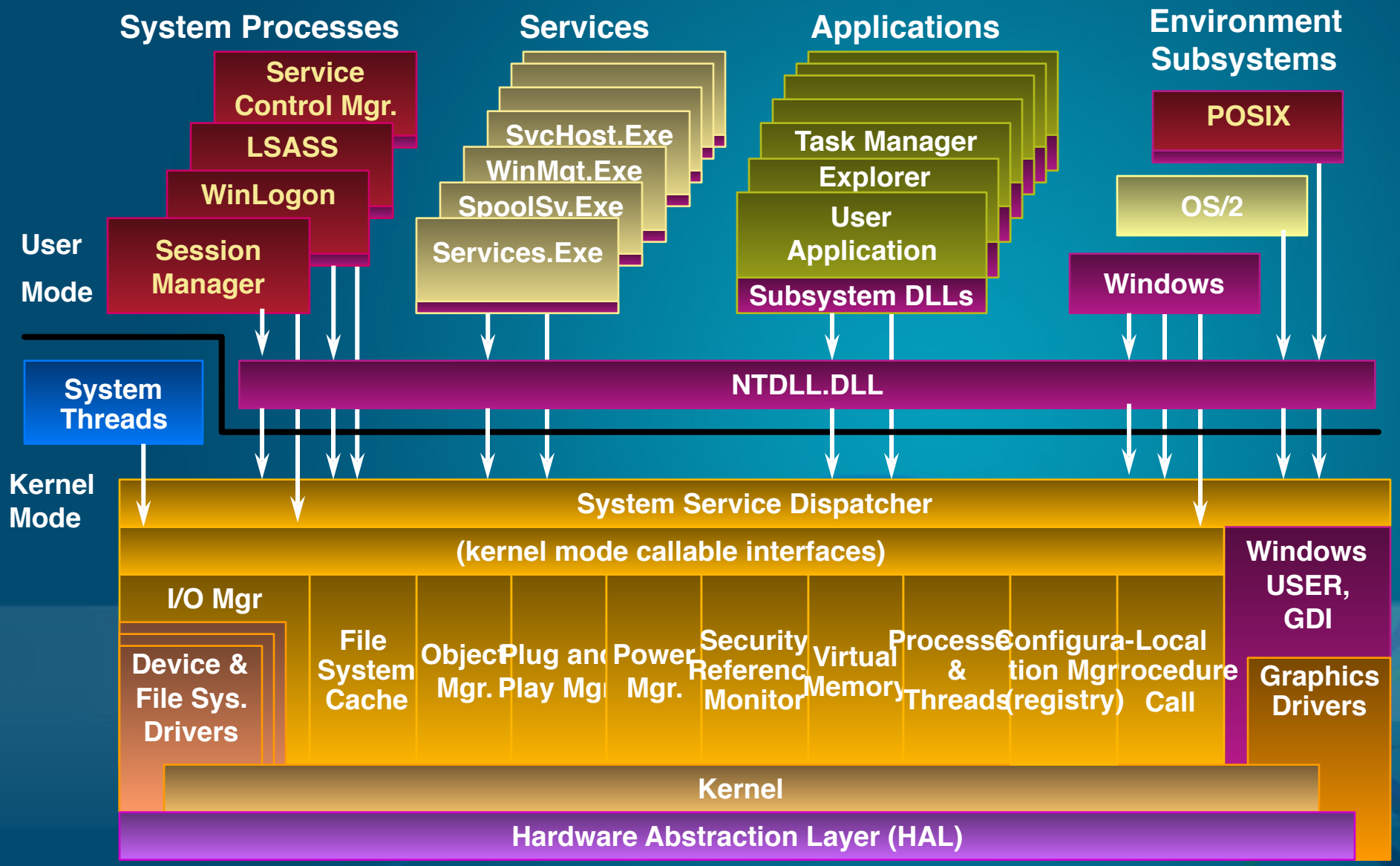
About the Speaker: Mark Russinovich

- Co-author of *Inside Windows 2000, 3rd Edition* and *Windows Internals, 4th edition* with David Solomon
- Senior Contributing Editor to Windows IT Pro Magazine
  - Co-authors Windows Power Tools column
- Author of tools on www.sysinternals.com
- Microsoft Most Valuable Professional (MVP)
- Co-founder and chief software architect of Winternals Software (www.winternals.com)
- Ph.D. in Computer Engineering

Purpose of Tutorial

- Give Windows developers a foundation understanding of the system's kernel architecture
  - Design better for performance & scalability
  - Debug problems more effectively
  - Understand system performance issues
- We're covering a small, but important set of core topics:
  - The "plumbing in the boiler room"

System Architecture



Tools Used To Dig In

- Many tools available to dig into Windows OS internals without requiring source code
  - Helps to see internals behavior “in action”
  - Many of these tools are used in labs in the video and the book
- Several sources of tools
  - Support Tools (on Windows OS CD-ROM in \support\tools)
  - Resource Kit Tools
  - Sysinternals tools (www.sysinternals.com)
  - Windows Debugging Tools

Live Kernel Debugging

- Useful for investigating internal system state not available from other tools
  - Previously, required 2 computers (host and target)
  - Target would be halted while host debugger in use
- XP & later supports live local kernel debugging
  - Technically requires system to be booted / DEBUG to work correctly
  - But, not all commands work

LiveKD

- LiveKd makes more commands work on a live system
  - Works on NT4, Windows 2000, Windows XP, Server 2003, and Vista
  - Was originally shipped on *Inside Windows 2000* book CD-ROM – now is free on Sysinternals
  - Tricks standard Microsoft kernel debuggers into thinking they are looking at a crash dump
  - Does not guarantee consistent view of system memory
    - Thus can loop or fail with access violation
    - Just quit and restart

Outline

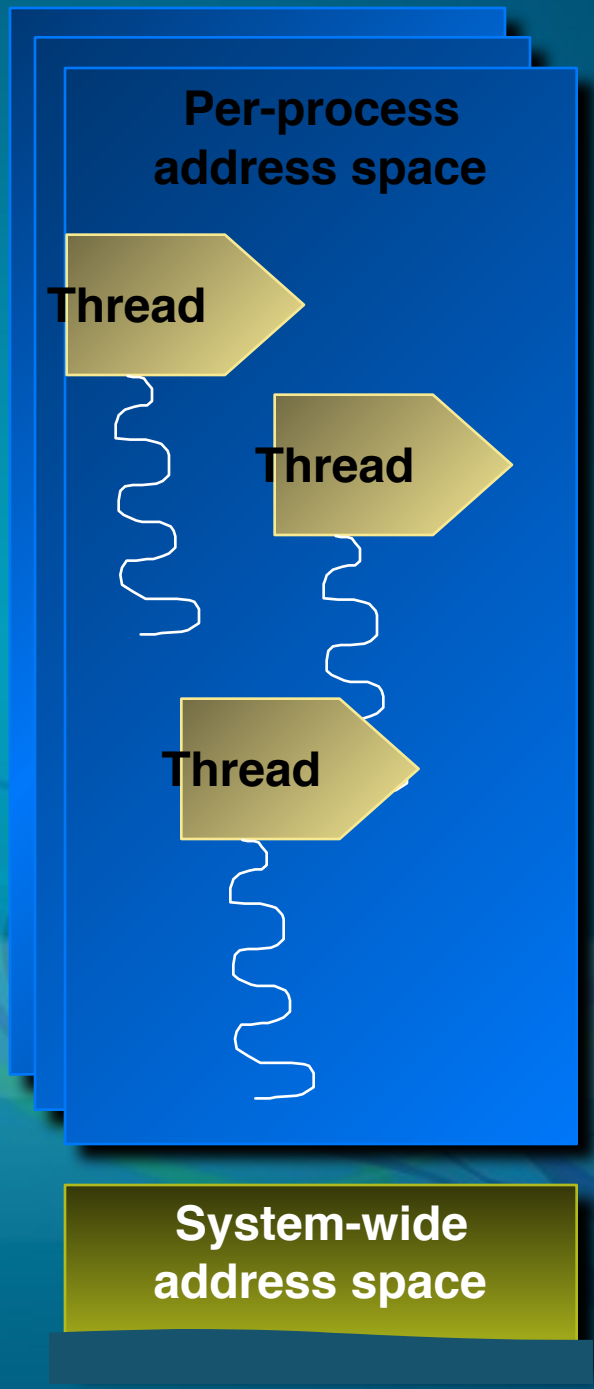
1. System Architecture
2. Processes and Thread Internals
3. Memory Management Internals
4. Security Internals

System Architecture

- Process Execution Environment
- Kernel Architecture
- Interrupt Handling
- Object Manager
- System Threads
- Process-based code
- Summary

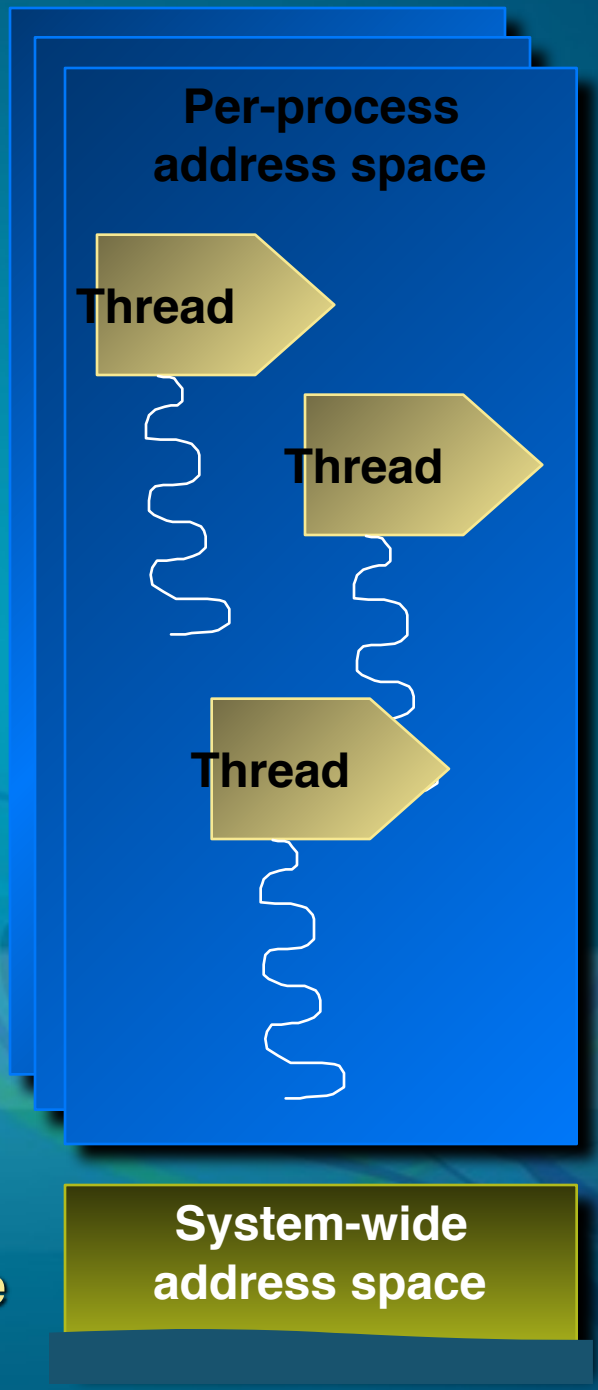
Processes And Threads

- What is a process?
  - Represents an instance of a running program
    - You create a process to run a program
    - Starting an application creates a process
  - Process defined by
    - Address space
    - Resources (e.g., open handles)
    - Security profile (token)
- System call
  - Primary argument to CreateProcess is image file name (or command line)



Processes And Threads

- What is a thread?
  - An execution context within a process
  - Unit of scheduling (threads run, processes don't run)
  - All threads in a process share the same per-process address space
    - Services provided so that threads can synchronize access to shared resources (critical sections, mutexes, events, semaphores)
  - All threads in the system are scheduled as peers to all others, without regard to their "parent" process
- System call:
  - Primary argument to CreateThread is a function entry point address
- Linux:
  - No threads per-se
  - Tasks can act like Windows threads by sharing handle table, PID and address space

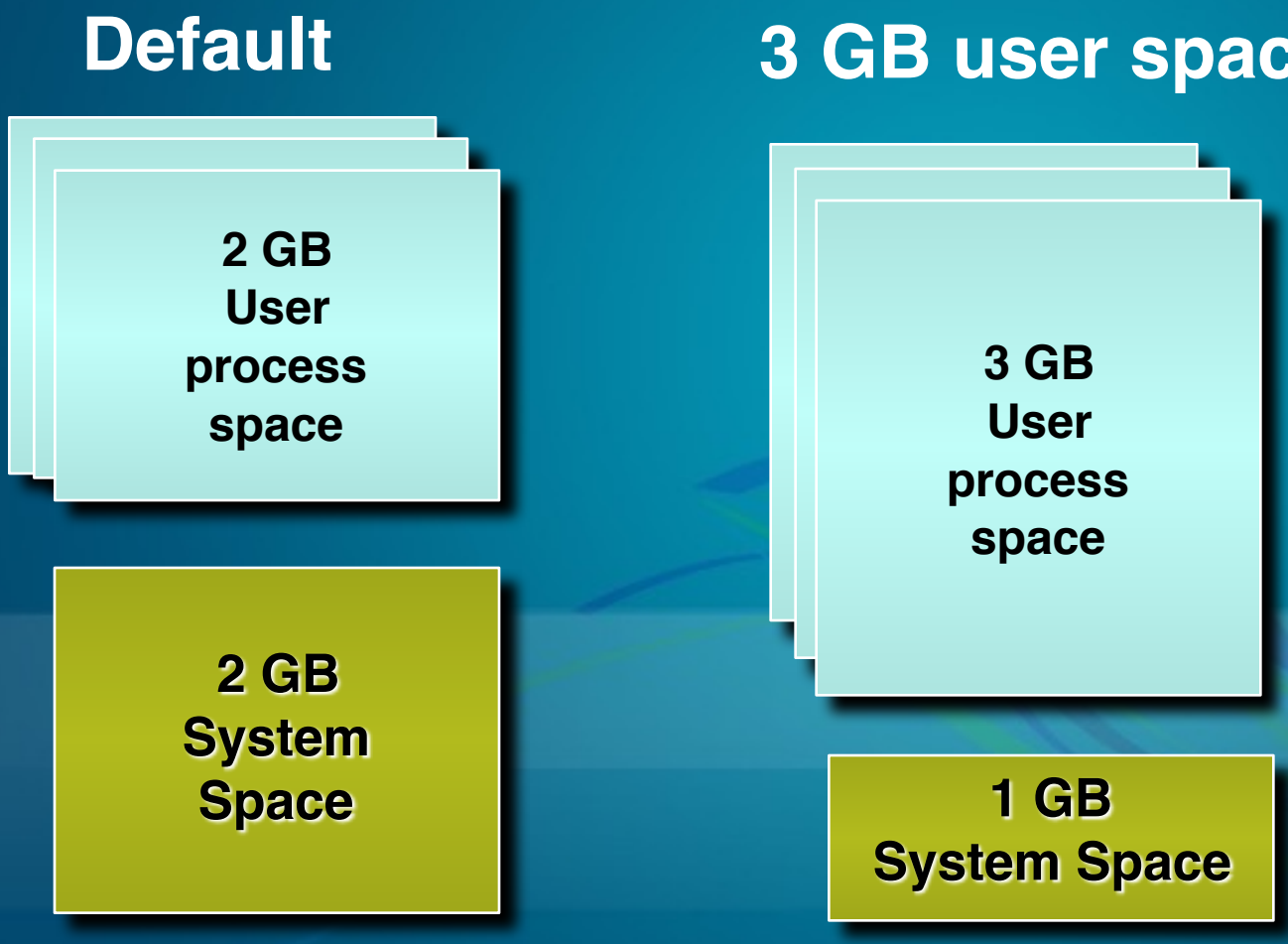


Processes And Threads

- Every process starts with one thread
  - First thread executes the program's "main" function
    - Can create other threads in the same process
    - Can create additional processes
- Why divide an application into multiple threads?
  - Perceived user responsiveness, parallel/background execution
    - Examples: Word background print – can continue to edit during print
  - Take advantage of multiple processors
    - On an MP system with n CPUs, n threads can literally run at the same time
    - Question: Given a single threaded application, will adding a second processor make it run faster?
  - Does add complexity
    - Synchronization
    - Scalability well is a different question…
      - Number of multiple runnable threads versus number CPUs
      - Having too many runnable threads causes excess context switching

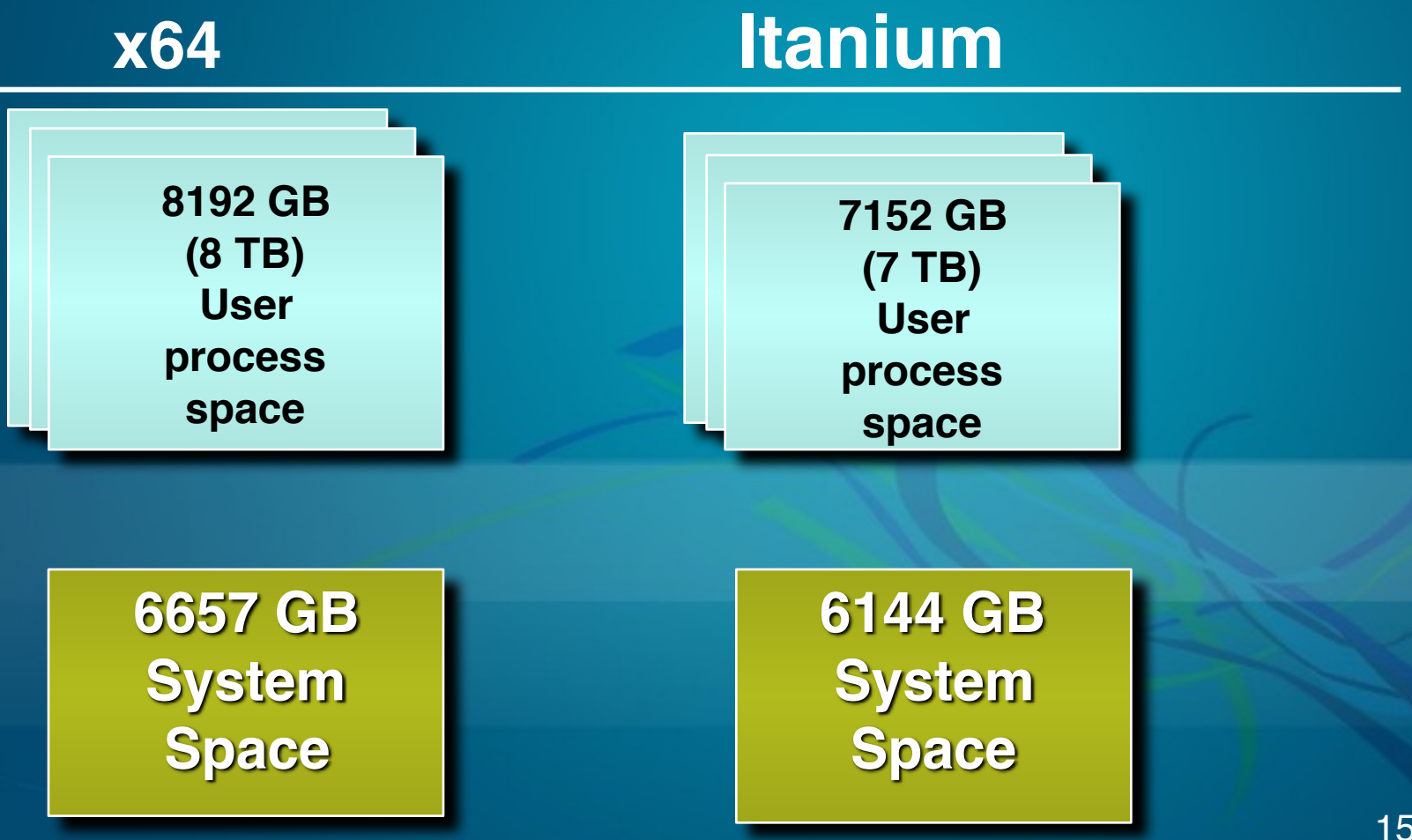
32-bit x86 Address Space

- 32-bits = 4 GB



64-bit Address Spaces

- 64-bits = 17,179,869,184 GB
  - x64 today supports 48 bits virtual = 262,144 GB
  - IA-64 today support 50 bits virtual = 1,048,576 GB

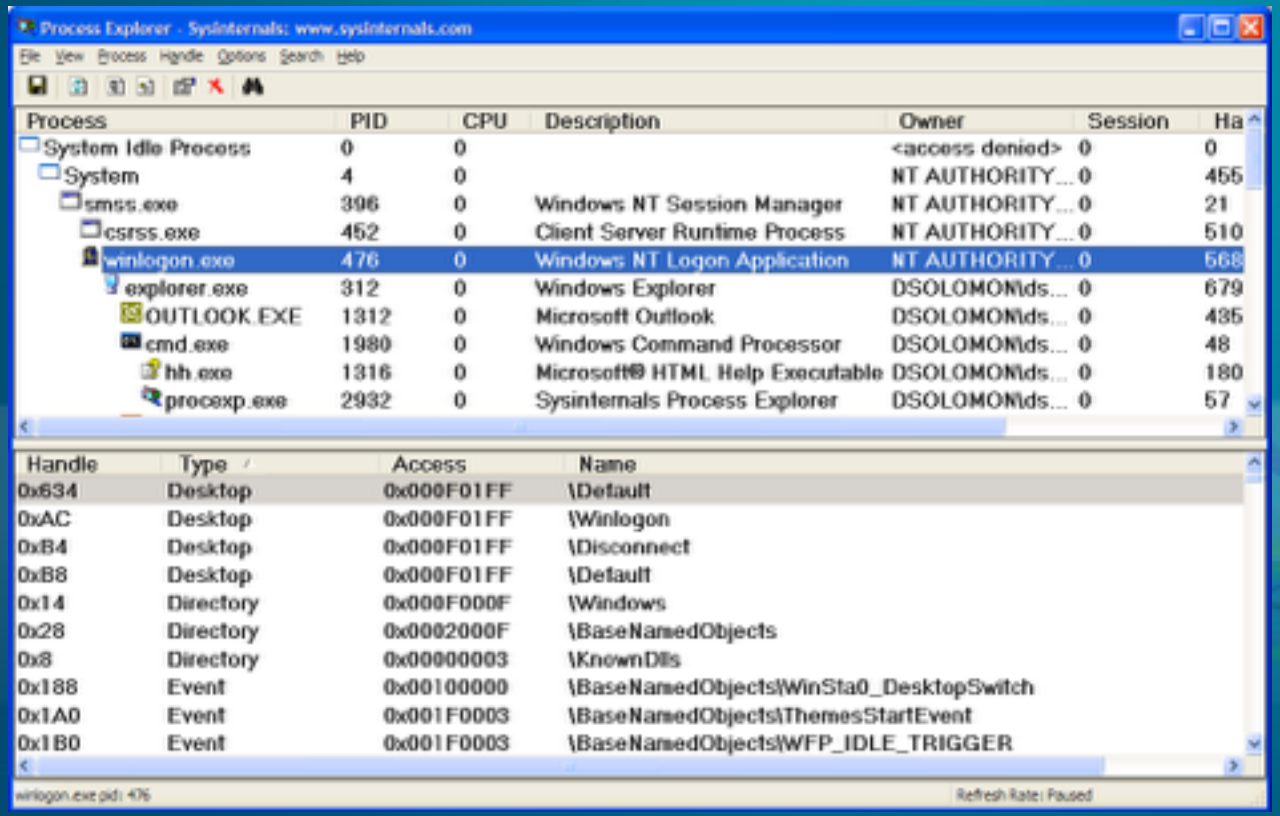


Memory Protection Model

- No user process can touch another user process address space (without first opening a handle to the process, which means passing through NT security)
  - Separate process page tables prevent this
  - “Current” page table changed on context switch from a thread in 1 process to a thread in another process
- No user process can touch kernel memory
  - Page protection in process page tables prevent this
  - OS pages only accessible from “kernel mode”
    - x86: Ring 0, Itanium: Privilege Level 0
  - Threads change from user to kernel mode and back (via a secure interface) to execute kernel code
    - Does not affect scheduling (not a context switch)

Process Explorer (Sysinternals)

- "Super Task Manager"
  - Shows full image path, command line, environment variables, parent process, thread details, security access token, open handles, loaded DLLs & mapped files



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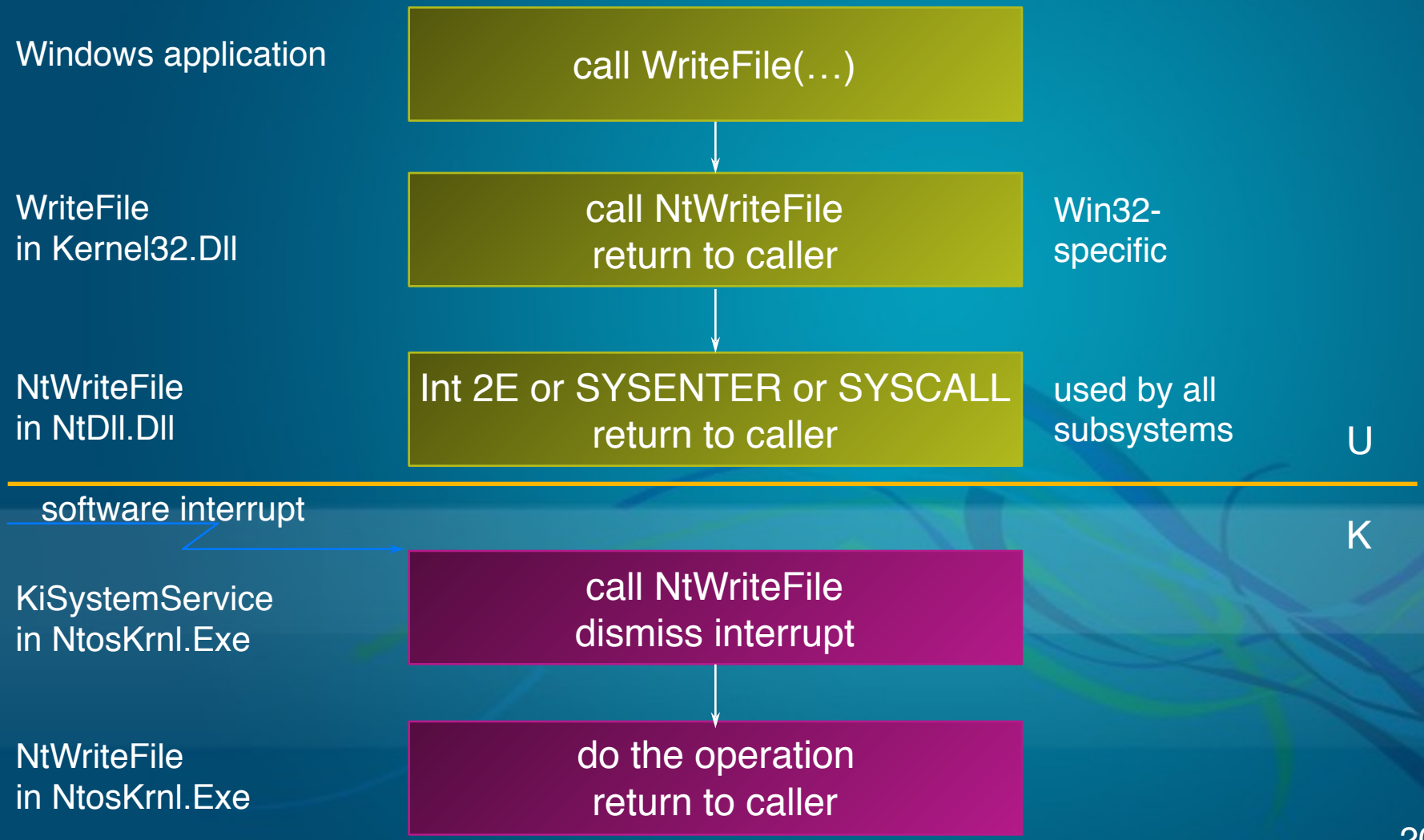
Windows Kernel Evolution

- Basic kernel architecture has remained stable while system has evolved
  - Windows 2000: major changes in I/O subsystem (plug & play, power management, WDM), but rest similar to NT4
  - Windows XP & Server 2003: modest upgrades as compared to the changes from NT4 to Windows 2000
- Internal version numbers confirm this:
  - Windows 2000 was 5.0
  - Windows XP is 5.1
  - Windows Server 2003 is 5.2
  - Windows Vista is 6.0

Example
Invoking a Win32 Kernel API



NTOSKRNL.EXE

- Core operating system image
  - Contains Executive and Kernel
- Four retail variations:

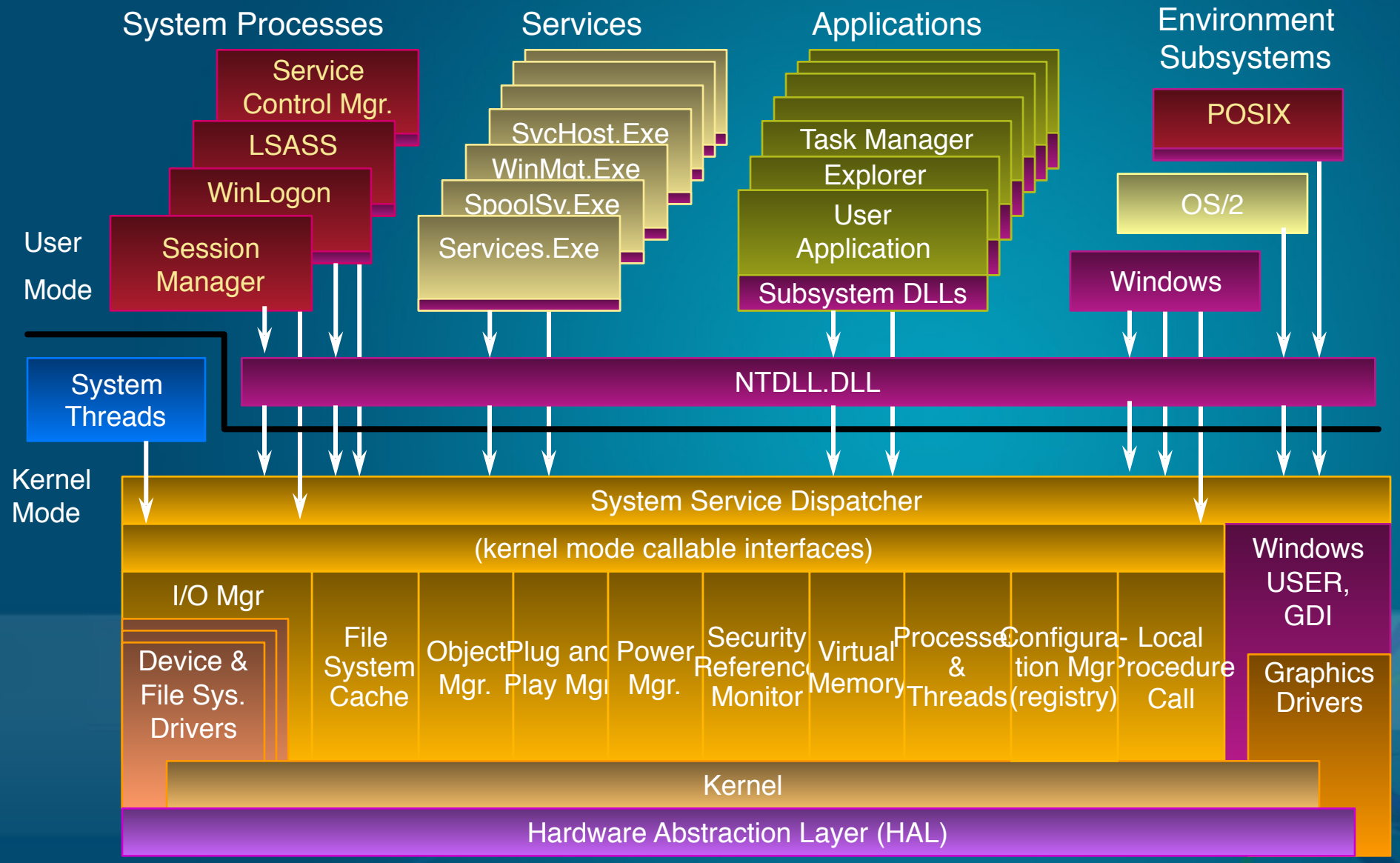
NTOSKRNL.EXE	Uniprocessor
NTKRNLMP.EXE	Multiprocessor

32-bit Windows PAE versions (for DEP & >4GB RAM):

NTKRNLPA.EXE	Uniprocessor w/extended addressing support
NTKRPAMP.EXE	Multiprocessor w/extended addressing support

- Vista: no uniprocessor kernel

System Architecture



Executive

- Upper layer of the operating system
- Provides “generic operating system” functions (“services”)
  - Process Manager
  - Object Manager
  - Cache Manager
  - LPC (local procedure call) Facility
  - Configuration Manager
  - Memory Manager
  - Security Reference Monitor
  - I/O Manager
  - Power Manager
  - Plug-and-Play Manager
- Almost completely portable C code
- Runs in kernel (“privileged”, ring 0) mode
- Most interfaces to executive services not documented

Kernel

- Lower layers of the operating system
  - Implements processor-dependent functions (x86 versus Itanium, etc.)
  - Also implements many processor-independent functions that are closely associated with processor-dependent functions
- Main services
  - Thread waiting, scheduling, and context switching
  - Exception and interrupt dispatching
  - Operating system synchronization primitives (different for MP versus UP)
  - A few of these are exposed to user mode
- Not a classic “microkernel”
  - shares address space with rest of kernel-mode components

HAL – Hardware Abstraction Layer

- Responsible for a small part of “hardware abstraction”
  - Components on the motherboard not handled by drivers
    - System timers, Cache coherency, and flushing
    - SMP support, Hardware interrupt priorities
- Subroutine library for the kernel and device drivers
  - Isolates OS & drivers from platform-specific details
  - Presents uniform model of I/O hardware interface to drivers
- Reduced role in Windows 2000
  - Bus support moved to bus drivers
  - Majority of HALs are vendor-independent

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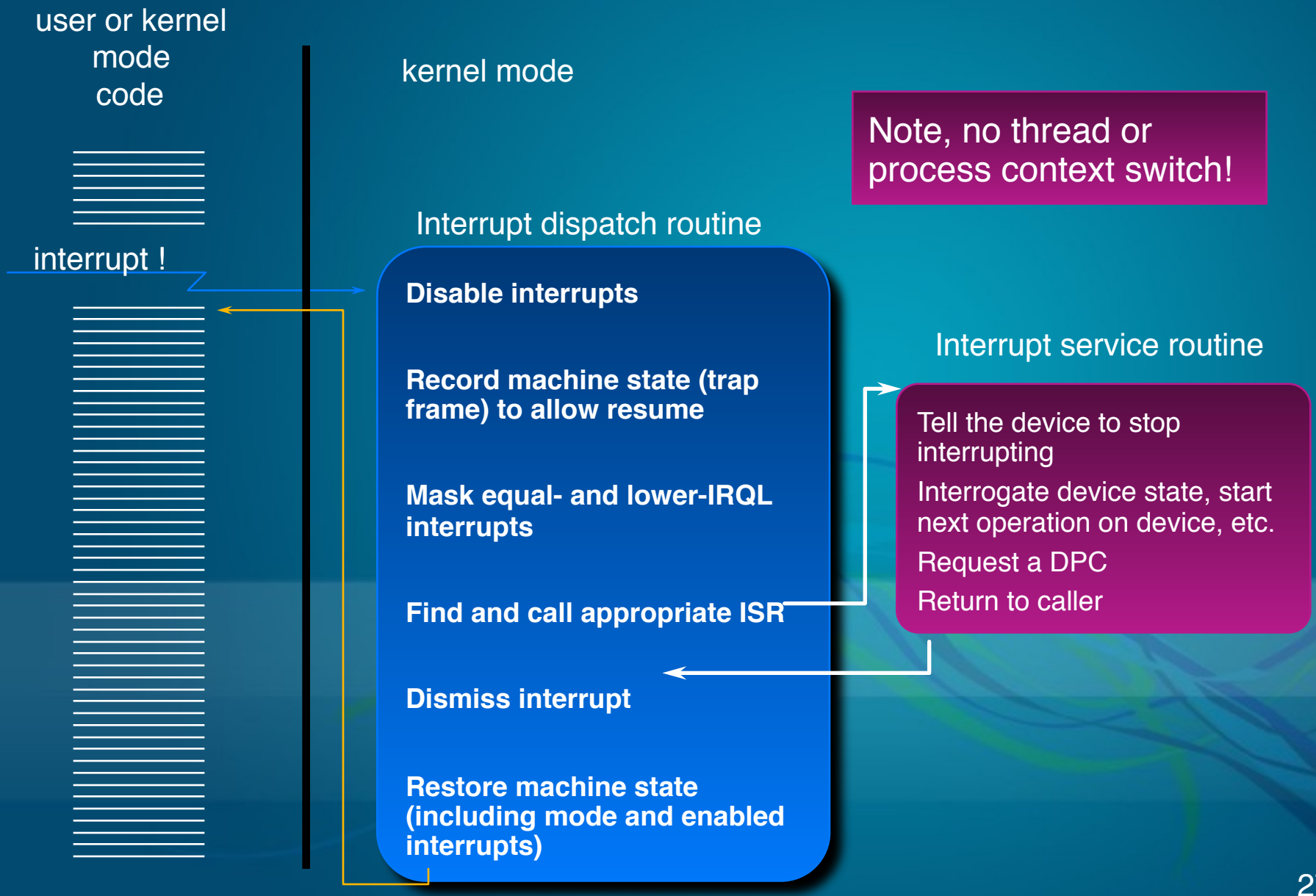
Interrupt Dispatching



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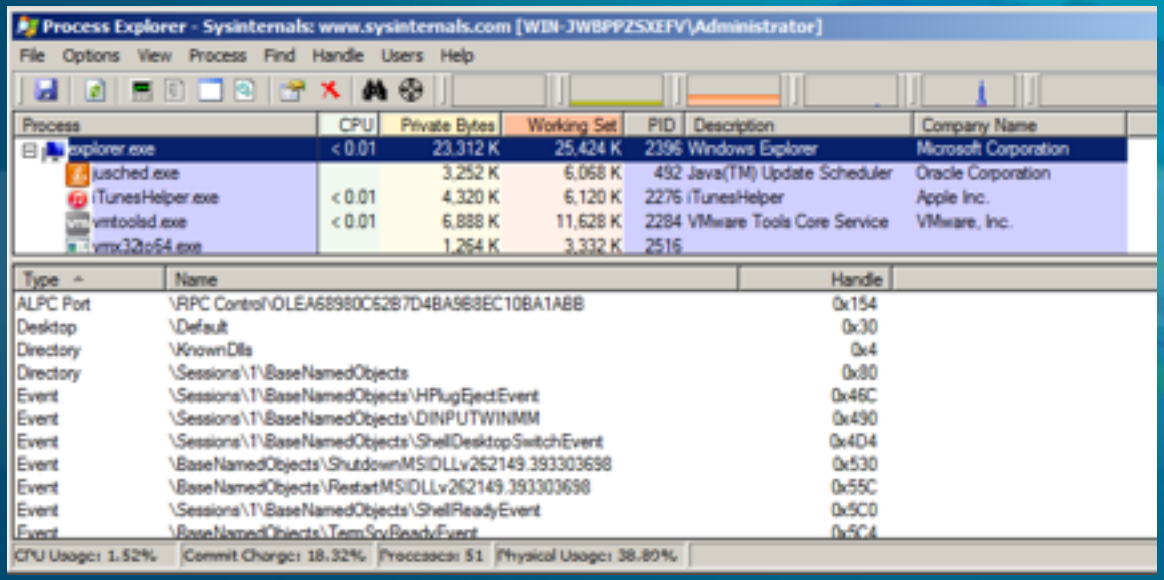
Handles And Security

- Process handle table
  - Is unique for each process
  - But is in system address space, hence cannot be modified from user mode
  - Hence, is trusted
- Security checks are made when handle table entry is created
  - i.e. at CreateXxx time
  - Handle table entry indicates the “validated” access rights to the object
    - Read, Write, Delete, Terminate, etc.
  - No need to revalidate on each request

Examining Open Handles: Sysinternals Tools

- Process Explorer
- View, Lower Pane View, Handles
- Right-click column header, select column "Handle Value"



Viewing Open Handles

- Handle View
  - By default, shows named objects
    - Click on Options->Show Unnamed Objects
- Uses:
  - Solve file locked errors
    - Can search to determine what process is holding a file or directory open
    - Can even close an open files (be careful!)
  - Understand resources used by an application
  - Detect handle leaks using refresh difference highlighting
  - View the state of synchronization objects (mutexes, semaphores, events)

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System Threads

- Functions in OS and some drivers that need to run as real threads
  - E.g., need to run concurrently with other system activity, wait on timers, perform background “housekeeping” work
  - Always run in kernel mode
  - Not non-preemptible (unless they raise IRQL to 2 or above)
  - For details, see DDK documentation on PsCreateSystemThread
- What process do they appear in?
  - “System” process (Windows NT 4.0: PID 2, Windows 2000: PID 8, Windows XP: PID 4)
  - In Windows 2000 and later, windowing system threads (from Win32k.sys) appear in “csrss.exe” (Windows subsystem process)

Examples Of System Threads

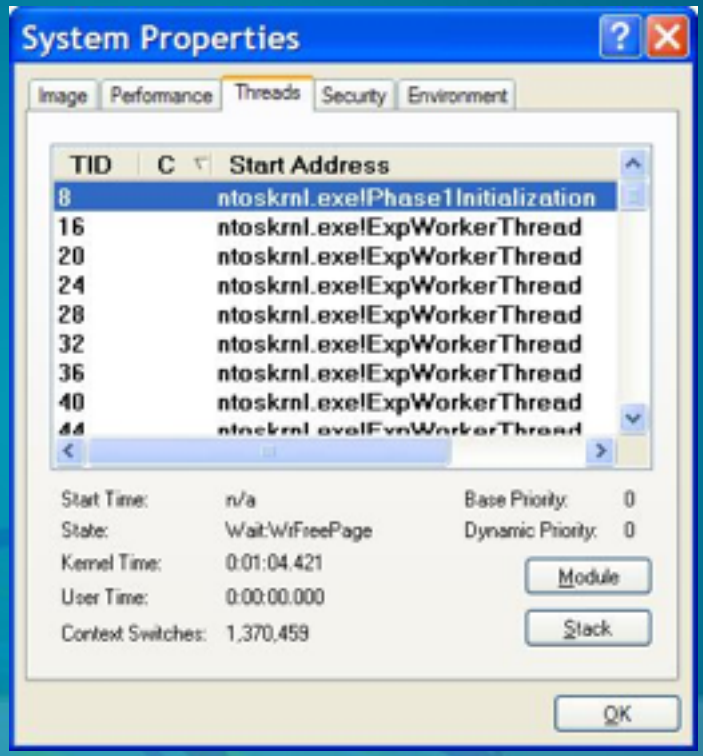
- Memory Manager
  - Modified Page Writer for mapped files
  - Modified Page Writer for paging files
  - Balance Set Manager
  - Swapper (kernel stack, working sets)
  - Zero page thread (thread 0, priority 0)
- Security Reference Monitor
  - Command Server Thread
- Network
  - Redirector and Server Worker Threads
- Threads created by drivers for their exclusive use
  - Examples: Floppy driver, parallel port driver
- Pool of Executive Worker Threads
  - Used by drivers, file systems, …
  - Accessed via ExQueueWorkItem

Identifying System Threads

- If System threads are consuming CPU time, need to find out what code is running, since it could be any one of a variety of components
  - Pieces of OS (Ntoskrnl.exe)
  - File server worker threads (Srv.sys)
  - Other drivers
- To really understand what's going on, must find which driver a thread "belongs to"

Identifiying System Threads

- Process Explorer:
  - Double click on System process
  - Go to Threads tab and sort by CPU
    - To view call stack, must use kernel debugger
- Note: several threads run between clock ticks (or at high IRQL) and thus don't appear to run
  - Watch context switch count



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Process-Based Code

- OS components that run in separate executables (.exes), in their own processes
  - Started by system
  - Not tied to a user logon
- Three types
  - Environment subsystems (already described)
  - System startup processes
    - Note: "system startup processes" is not an official Microsoft defined name
  - Windows Services
- Let's examine the system process "tree"
  - Use Tlist /T or Process Explorer

Process-Based NT Code
System Startup Processes

- First two processes aren't real processes
  - Not running a user mode .EXE
  - No user-mode address space
  - Different utilities report them with different names
  - Data structures for these processes (and their initial threads) are "pre-created" in NtosKrnl.Exe and loaded along with the code

(Idle)	Process id 0
	Part of the loaded system image
	Home for idle thread(s) (not a real process nor real threads)
	Called "System Process" in many displays
(System)	Process id 2 (8 in Windows 2000; 4 in XP)
	Part of the loaded system image
	Home for kernel-defined threads (not a real process)
	Thread 0 (routine name Phase1Initialization) launches the first "real" process, running smss.exe...
	...and then becomes the zero page thread

Process-Based NT Code
System Startup Processes

smss.exe Session Manager
The first “created” process
Takes parameters from \HKEY_LOCAL_MACHINE\System\CurrentControlSet\Control\Session Manager
Launches required subsystems (csrss) and then winlogon

csrss.exe Windows subsystem

---

winlogon.exe Logon process: Launches services.exe & lsass.exe; presents first login prompt
When someone logs in, launches apps in \Software\Microsoft\Windows NT\WinLogon\Userinit

services.exe Service Controller; also, home for many NT-supplied services
Starts processes for services not part of services.exe (driven by \Registry\Machine\System\CurrentControlSet\Services)

lsass.exe Local Security Authentication Server

userinit.exe Started after logon; starts Explorer.exe (see \Software\Microsoft\Windows NT\CurrentVersion\WinLogon\Shell) and exits (hence Explorer appears to be an orphan)

explorer.exe and its children are the creators of all interactive apps

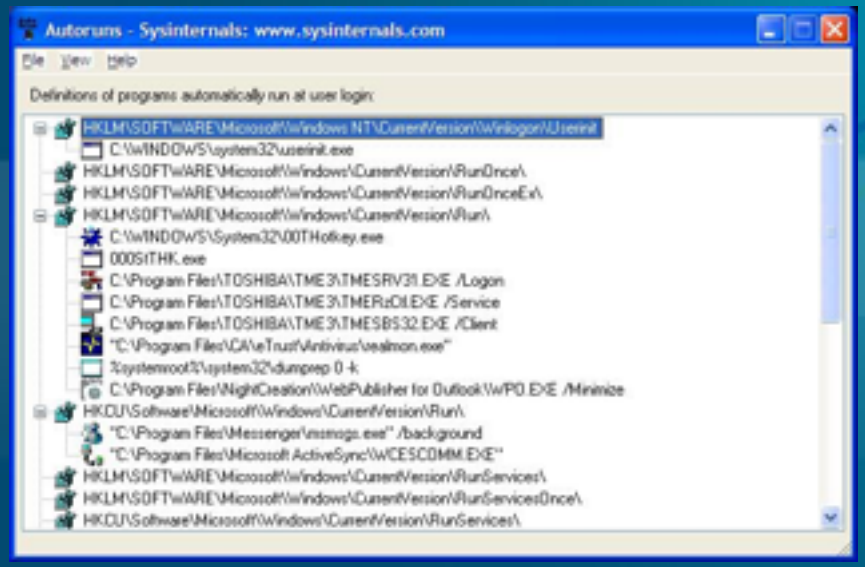
Logon Process

- Winlogon sends username/password to Lsass
  - Either on local system for local logon, or to Netlogon service on a domain
  - Windows XP enhancement: Winlogon doesn't wait for Workstation service to start if
    - Account doesn't depend on a roaming profile
    - Domain policy that affects logon hasn't changed since last logon
    - Controller for a network logon
- Creates a process to run
  HKLM\Software\Microsoft\Windows NT\CurrentVersion\WinLogon\Userinit
  - By default: Userinit.exe
  - Runs logon script, restores drive-letter mappings, starts shell
- Userinit creates a process to run
  HKLM\Software\Microsoft\Windows NT\CurrentVersion\WinLogon\Shell
  - By default: Explorer.exe
- There are other places in the Registry that control programs that start at logon

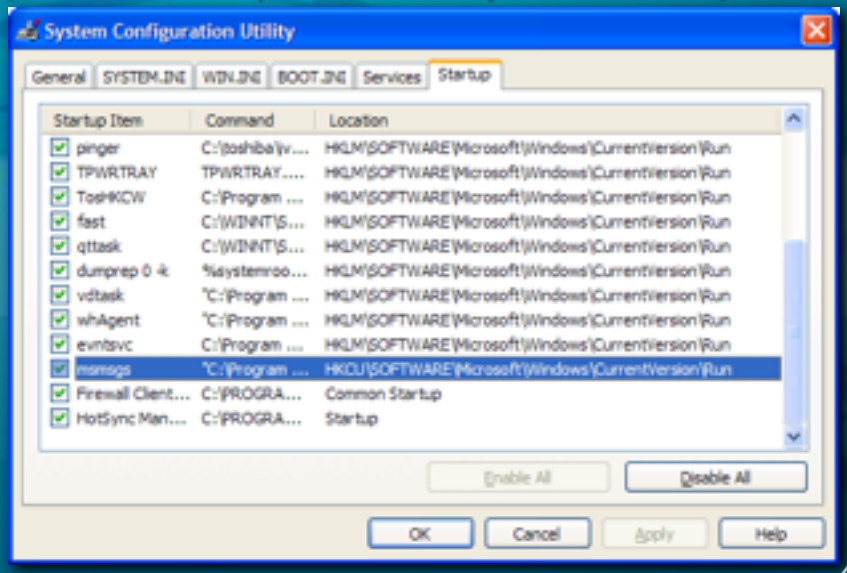
Processes Started at Logon

- Displays order of processes configured to start at log on time
- Also can use new XP built-in tool called "System Configuration Utility"
  - To run, click on Start->Help, then "Use Tools…", then System Configuration Utility
  - Only shows what's defined to start vs Autoruns which shows all places things CAN be defined to start

Autoruns (Sysinternals)



Msconfig
(in \Windows\pchealth\helpctr\binaries)

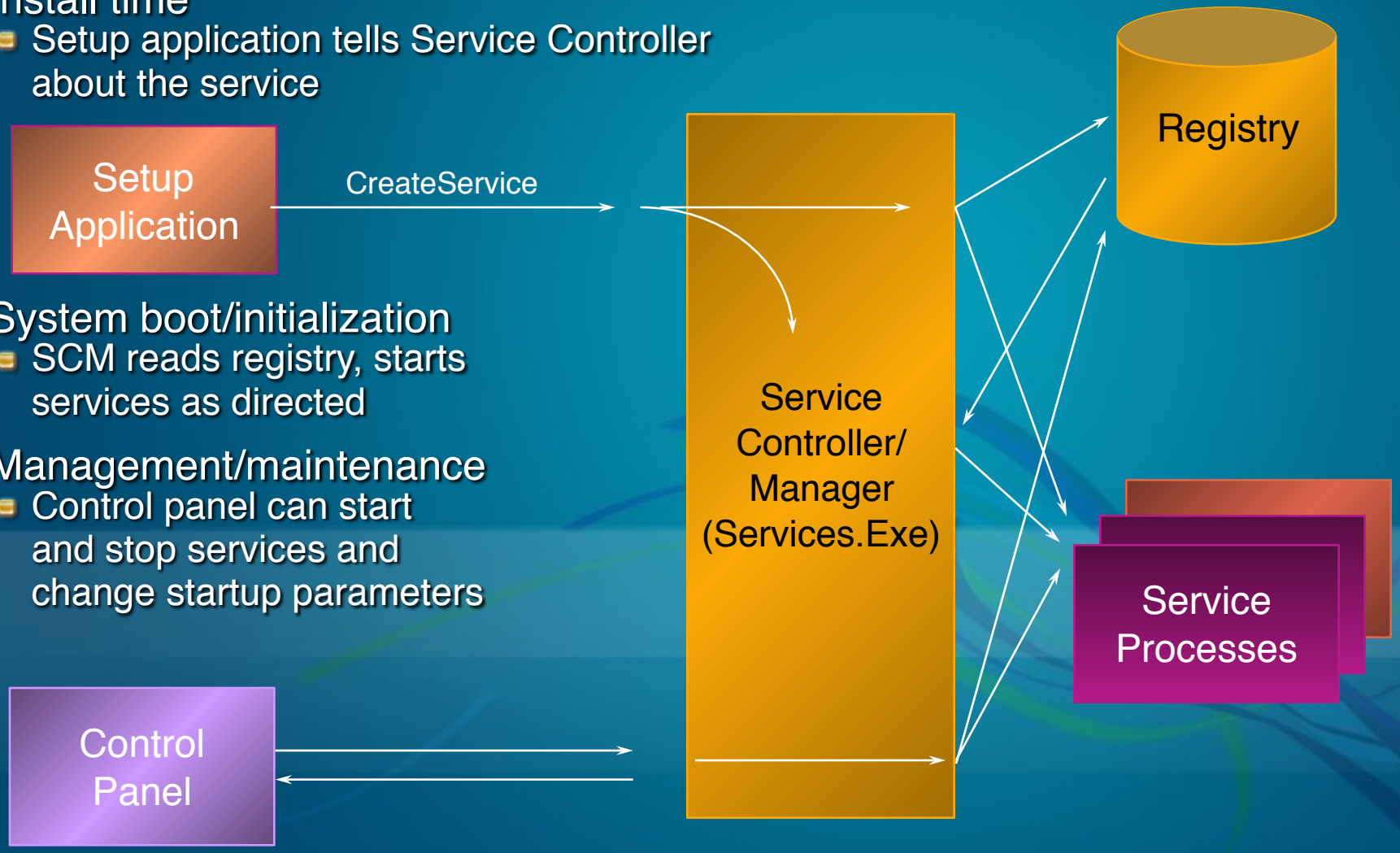


Windows Services

- An overloaded generic term
- A process created and managed by the Service Control Manager (Services.exe)
  - E.g. Solitaire can be configured as a service, but is killed shortly after starting
- Similar in concept to Unix daemon processes
  - Typically configured to start at boot time (if started while logged on, survive logoff)
  - Typically do not interact with the desktop
- Note: Prior to Windows 2000 this is one way to start a process on a remote machine (now you can do it with WMI)

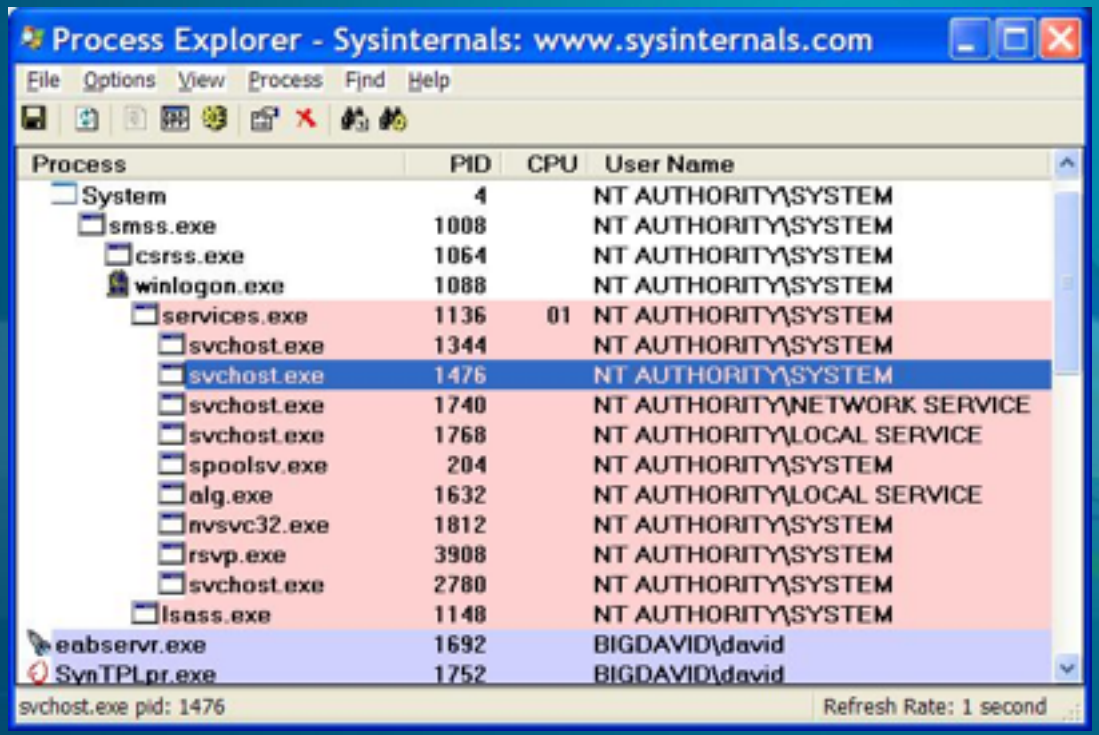
Life Of A Service

- Install time
  - Setup application tells Service Controller about the service
- System boot/initialization
  - SCM reads registry, starts services as directed
- Management/maintenance
  - Control panel can start and stop services and change startup parameters



Viewing Service Processes

- Process Explorer colors Services pink by default



Svchost Mechanism

- Windows 2000 introduced generic Svchost.exe
  - Groups services into fewer processes
    - Improves system startup time
    - Conserves system virtual memory
  - Not user-configurable as to which services go in which processes
  - 3rd parties cannot add services to Svchost.exe processes
- Windows XP/2003 have more Svchost processes due to two new less privileged accounts for built-in services
  - LOCAL SERVICE, NETWORK SERVICE
  - Less rights than SYSTEM account
    - Reduces possibility of damage if system compromised
- On XP/2003, four Svchost processes (at least):
  - SYSTEM, SYSTEM (2nd instance – for RPC), LOCAL SERVICE, NETWORK SERVICE

Mapping Services To Service Processes

- Tlist /S (Debugging Tools) or Tasklist /svc (XP/2003) list internal name of services inside service processes
- Process Explorer shows more: external display name and description

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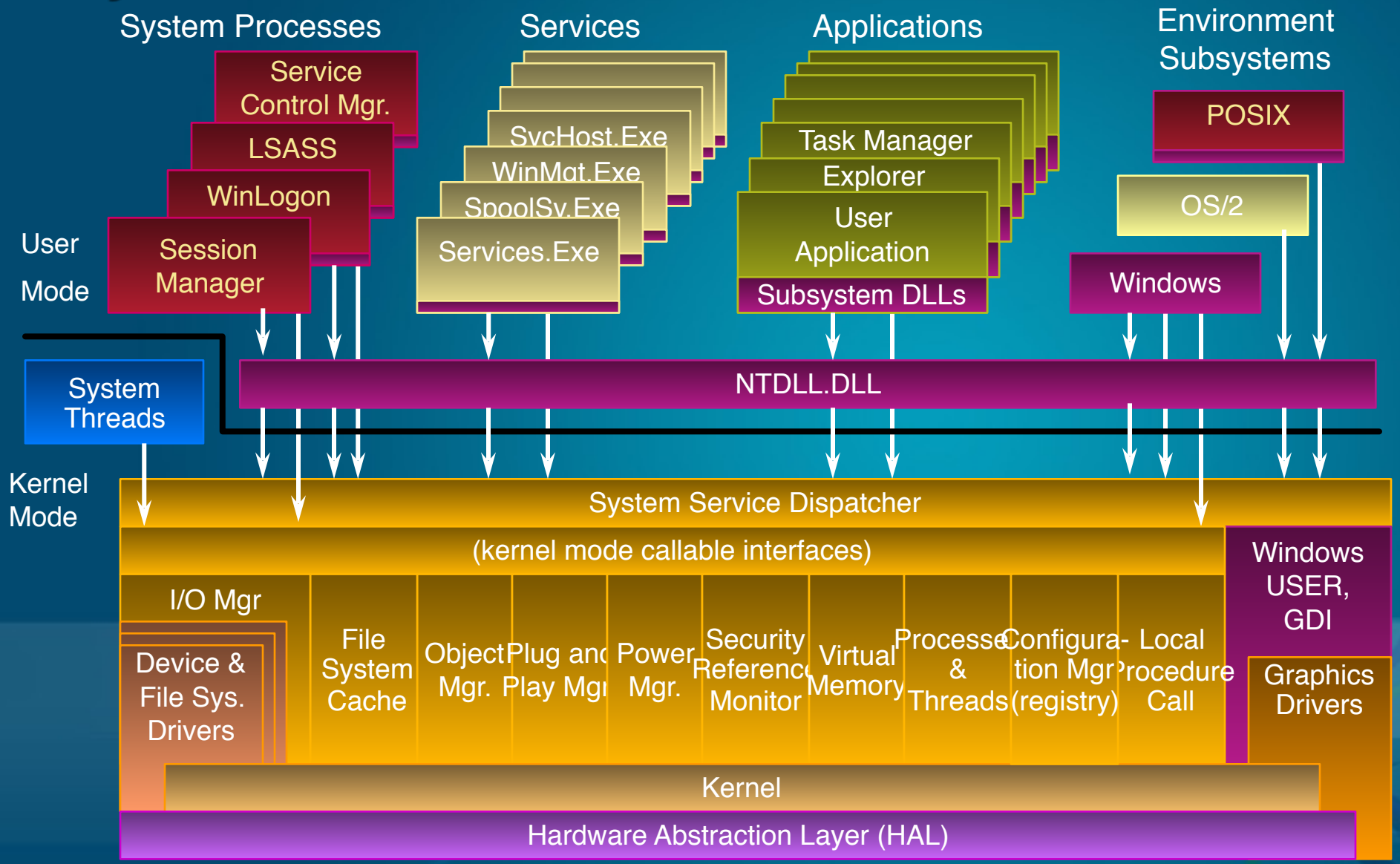
Four Contexts For Executing Code

- Full process and thread context
  - User applications
  - Windows Services
  - Environment subsystem processes
  - System startup processes
- Have thread context but no “real” process
  - Threads in “System” process
- Routines called by other threads/processes
  - Subsystem DLLs
  - Executive system services (NtReadFile, etc.)
  - GDI32 and User32 APIs implemented in Win32K.Sys (and graphics drivers)
- No process or thread context (“arbitrary thread context”)
  - Interrupt dispatching
  - Device drivers

System Architecture



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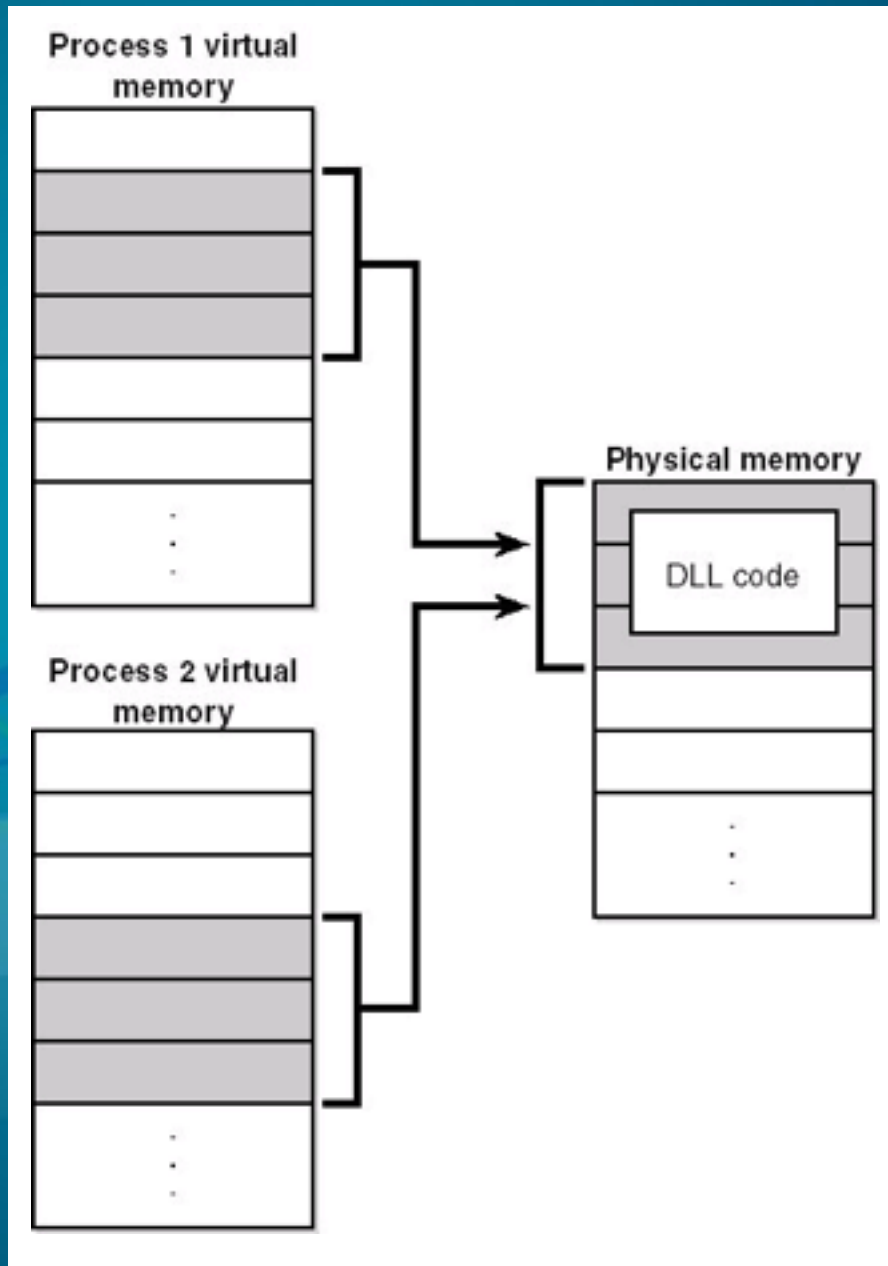
Memory Management

- Core Memory Management Services
- Working Set Management
- Unassigned Memory
- Page Files

Shared Memory

- Like most modern OSs, Windows provides a way for processes to share memory
  - High speed IPC (used by LPC, which is used by RPC)
  - Threads share address space, but applications may be divided into multiple processes for stability reasons
- Processes can also create shared memory sections
  - Called page file backed file mapping objects
  - Full Windows security
- It does this automatically for shareable pages
  - E.g., code pages in an .EXE



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Prefetch Mechanism

- File activity is traced and used to prefetch data the next time
  - First 10 seconds are monitored
    - Pages referenced & directories opened
  - Prefetch “trace file” stored in \Window\Prefetch
    - Name of .EXE-<hash of full path>.pf
- Also applies to system boot
  - First 2 minutes of boot process logged
    - Stops 30 seconds after the user starts the shell or 60 seconds after all services are started
  - Boot trace file: NTOSBOOT-B00DFAAD.pf

Prefetch Mechanism

- When application run again, system automatically
  - Reads in directories referenced
  - Reads in code and file data
    - Reads are asynchronous
    - But waits for all prefetch to complete
- In addition, every 3 days, system automatically defrags files involved in each application startup
- Bottom line: Reduces disk head seeks
  - This was seen to be the major factor in slow application/system startup

Memory Management

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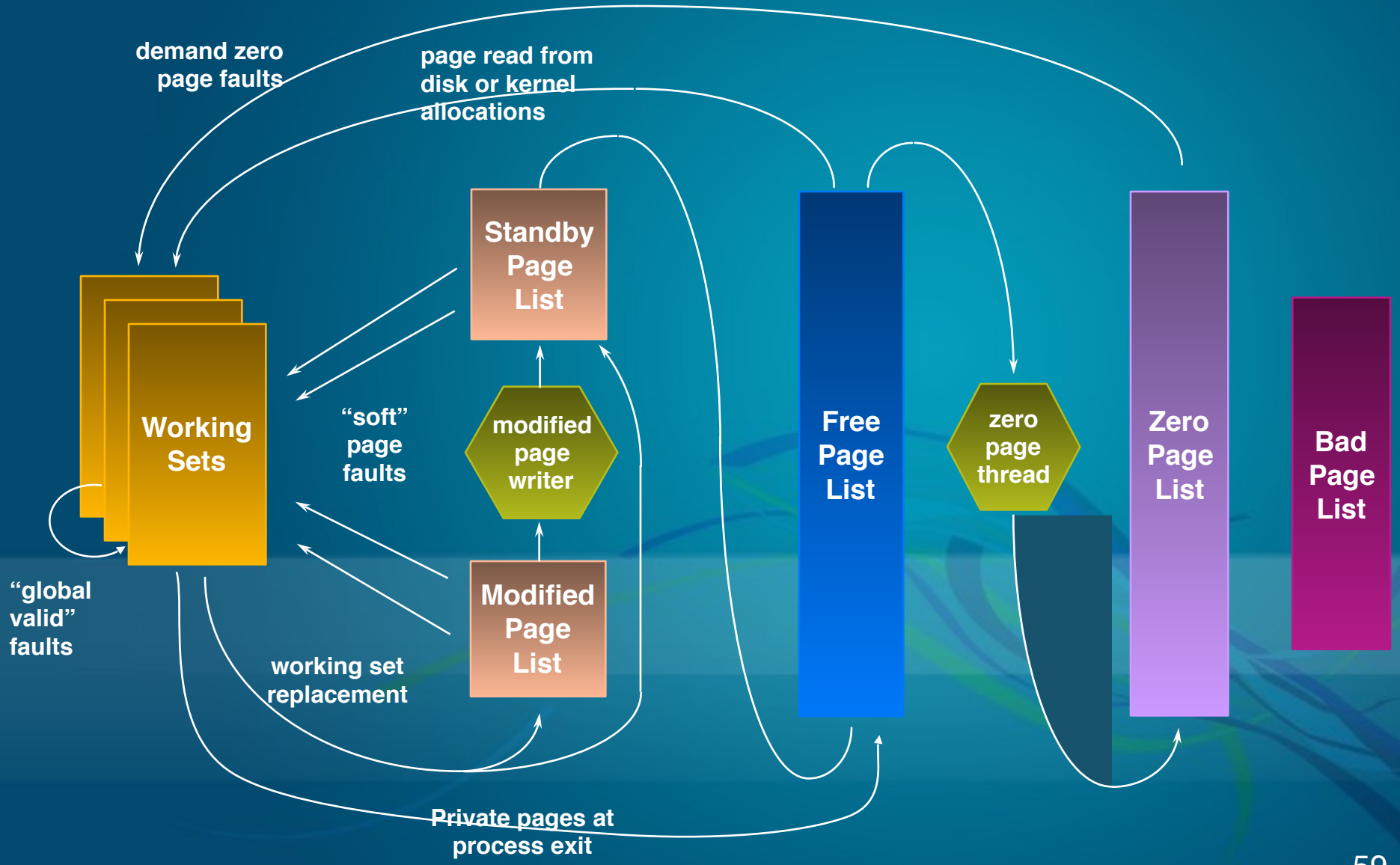
Managing Physical Memory

- System keeps unassigned physical pages on one of several lists
  - Free page list
  - Modified page list
  - Standby page list
  - Zero page list
  - Bad page list – pages that failed memory test at system startup
- Lists are implemented by entries in the “PFN database”
  - Maintained as FIFO lists or queues

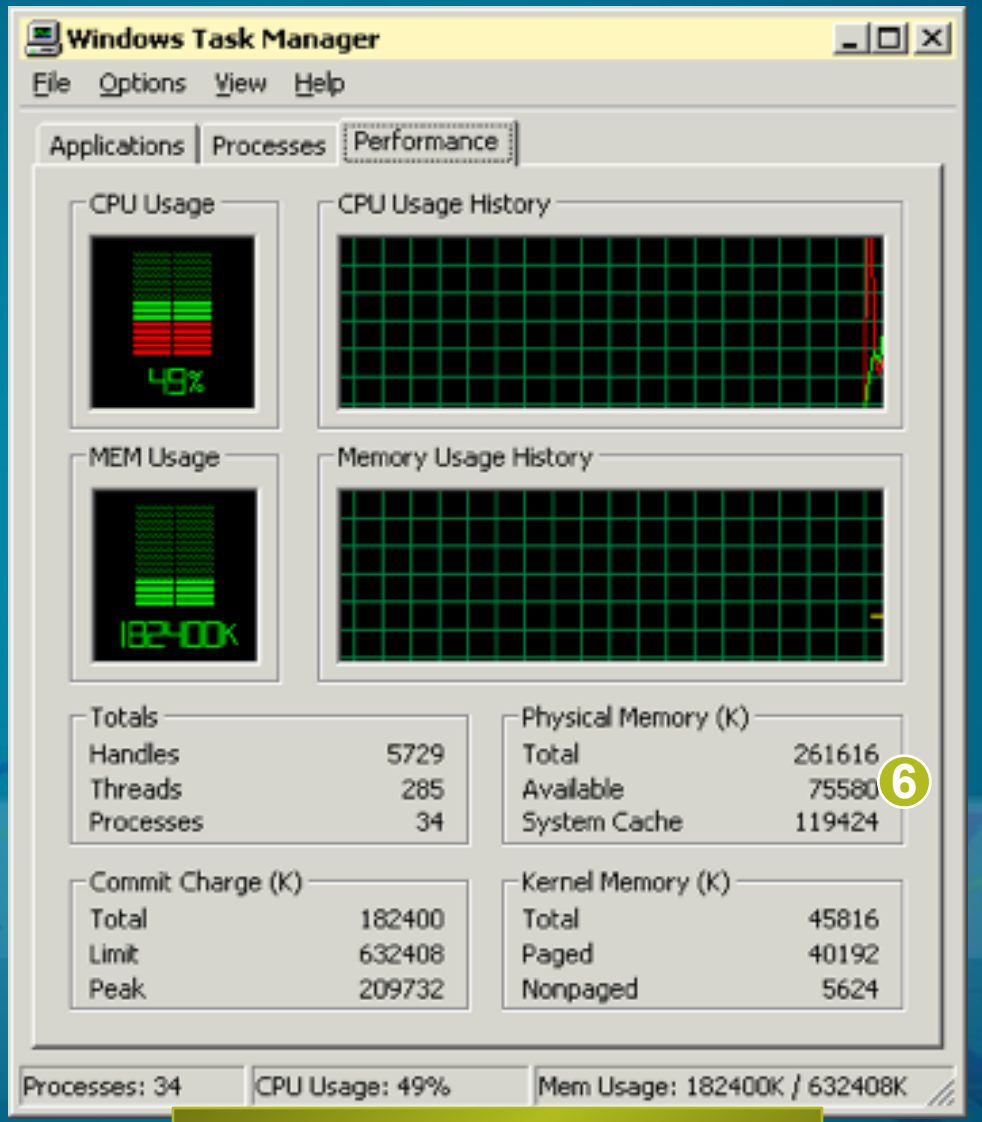
Paging Dynamics



Memory Management Information

Task Manager
Performance tab

- 6 "Available" = sum of free, standby, and zero page lists (physical)
- Majority are likely standby pages
- "System Cache" = size of standby list + size of system working set (file cache, paged pool, pageable OS/driver code & data)



Screen snapshot from:
Task Manager | Performance tab

Outline

1. System Architecture
2. Processes and Thread Internals
3. Memory Management Internals
4. Security Internals

Security

- Introduction
- Components
- Logon
- Protecting Objects
- Privileges

Windows Security Support

- Microsoft's goal was to achieve C2, which requires:
  - Secure Logon: NT provides this by requiring user name and password
  - Discretionary Access Control: fine grained protection over resources by user/group
  - Security Auditing: ability to save a trail of important security events, such as access or attempted access of a resource
  - Object reuse protection: must initialize physical resources that are reused e.g. memory, files
- Certifications achieved:
  - Windows NT 3.5 (workstation and server) with SP3 earned C2 in July 1995
  - In March 1999 Windows NT 4 with SP3 earned e3 rating from UK's Information Technology Security (ITSEC) – equivalent to C2
  - In November 1999 NT4 with SP6a earned C2 in stand-alone and networked environments

Windows Security Support

- Windows meets two B-level requirements:
  - Trusted Path Functionality: way to prevent trojan horses with “secure attention sequence” (SAS) - Ctrl-Alt-Del
  - Trusted Facility Management: ability to assign different roles to different accounts
    - Windows does this through account privileges (TBD later)

Common Criteria

- New standard, called Common Criteria (CC), is the new standard for software and OS ratings
  - Consortium of US, UK, Germany, France, Canada, and the Netherlands in 1996
  - Became ISO standard 15408 in 1999
  - For more information, see http://www.commoncriteriaportal.org/ and http://csrc.nist.gov/cc
- CC is more flexible than TCSEC trust ratings
  - Protection Profile collects security requirements
  - Security Target (ST) are security requirements that can be made by reference to a PP
- Windows 2000 was certified as compliant with the CC Controlled Access Protection Profile (CAPP) in October 2002
  - Windows XP and Server 2003 are undergoing evaluation

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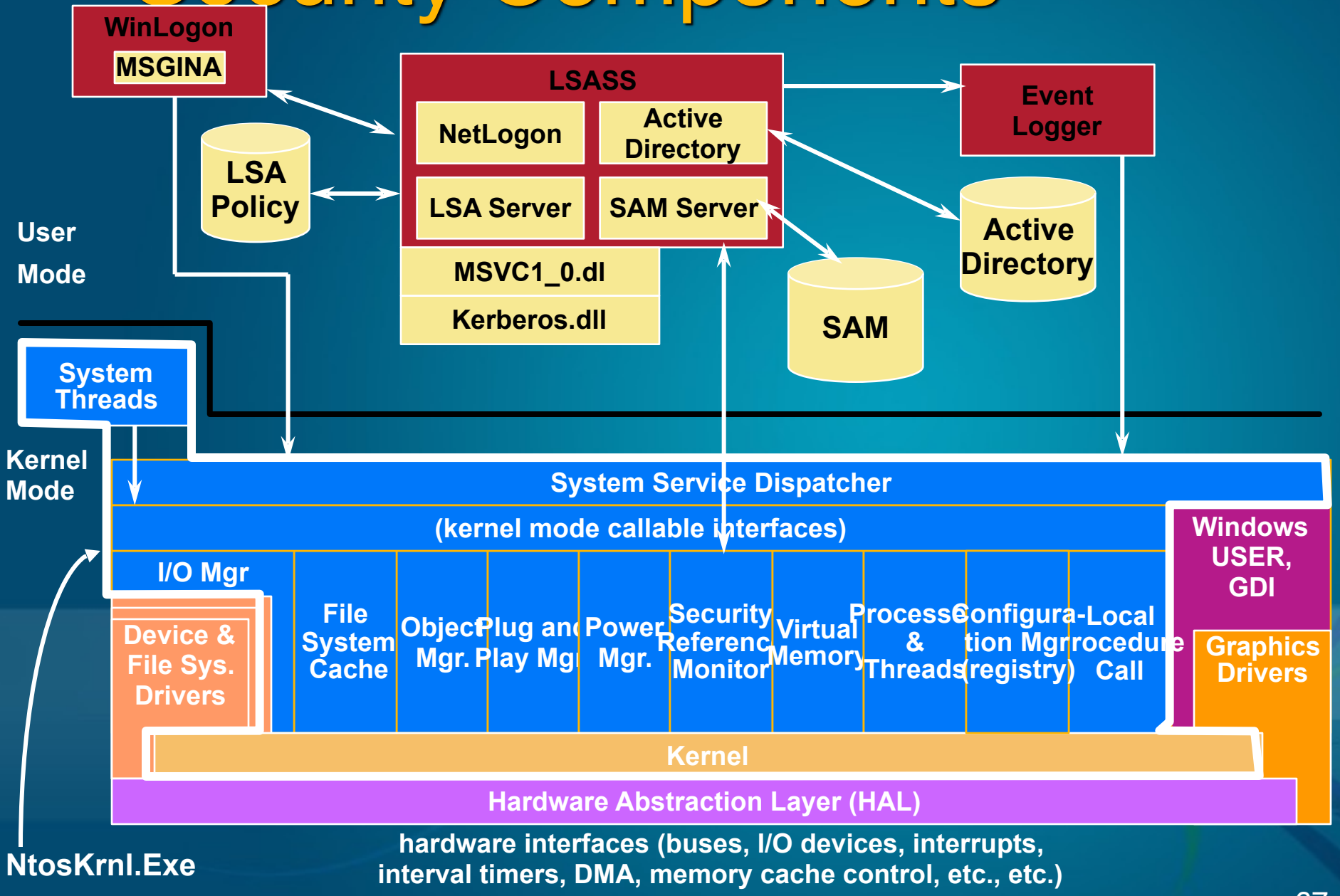
Security Components



Security Reference Monitor

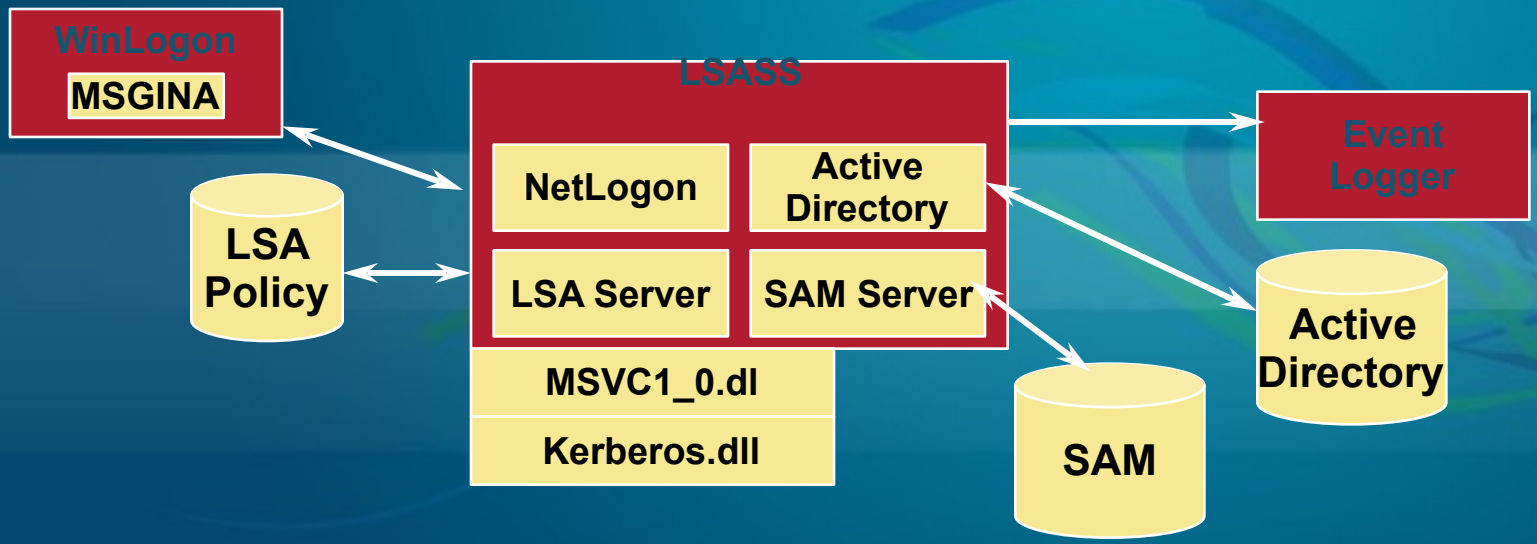
- Performs object access checks, manipulates privileges, and generates audit messages
- Group of functions in Ntoskrnl.exe
  - Some documented in DDK
  - Exposed to user mode by Windows API calls
- Demo: Open Ntoskrnl.exe with Dependency Walker and view functions starting with “Se”

Demo: Viewing Security Processes

- Run Process Explorer
- Collapse Explorer process tree and focus on upper half (system processes)

Security Components

- Local Security Authority
  - User-mode process (\Windows\System32\Lsass.exe) that implements policies (e.g. password, logon), authentication, and sending audit records to the security event log
  - LSASS policy database: registry key HKLM\SECURITY



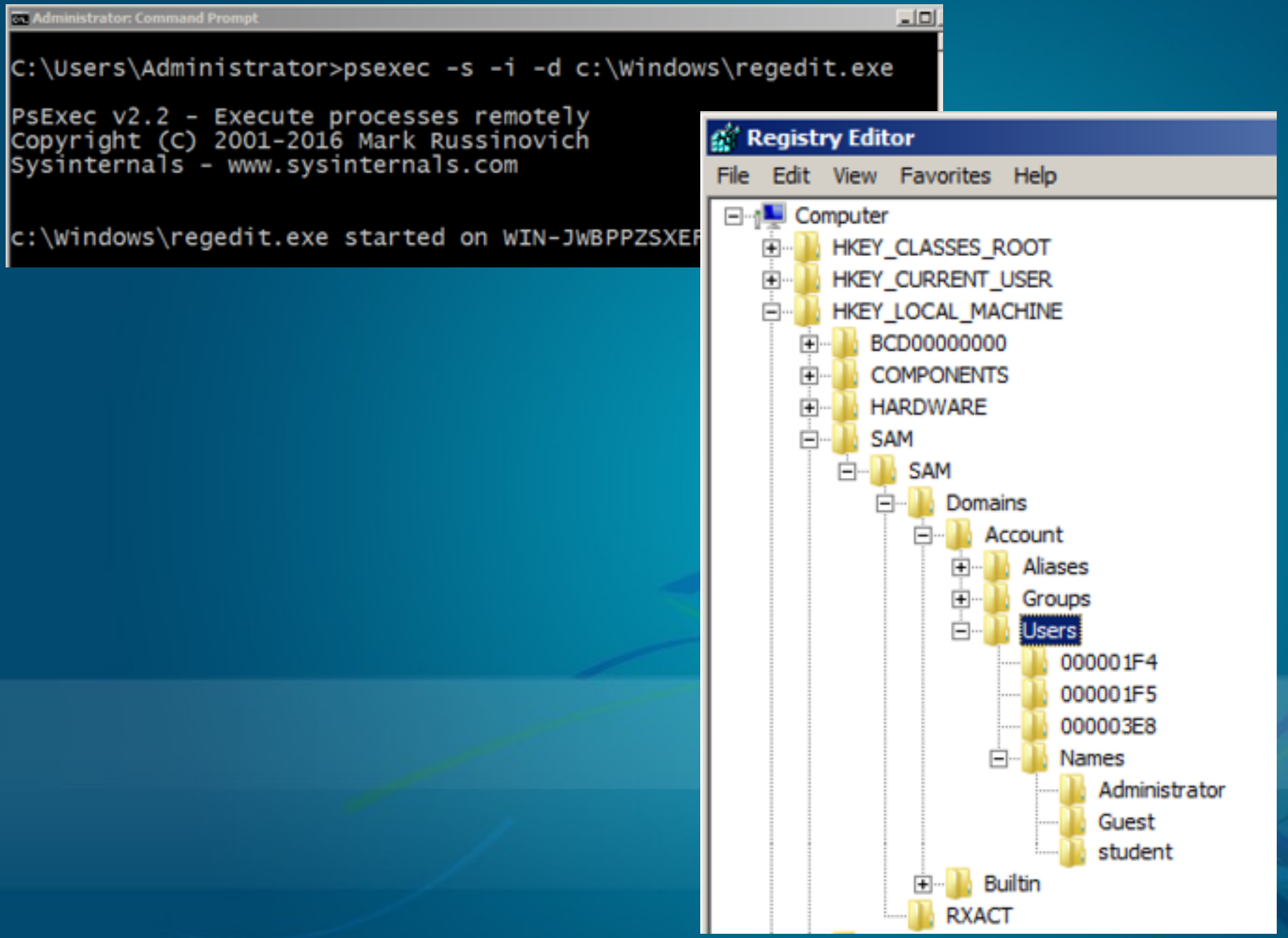
Demo: Looking at the SAM

- Look at HKLM\SAM permissions
  - SAM security allows only the local system account to access it
  - Run Regedit
  - Look at HKLM\SAM - nothing there?
  - Check permissions (right click->Permissions)
  - Close Regedit
- Look in HKLM\SAM
  - Running Regedit in the local system account allows you to view the SAM:

    psexec –s –i –d c:\windows\regedit.exe

    or

    sc create cmdassystem type= own type= interact
    binpath= "cmd /c start cmd /k“
    sc start cmdassystem
  - View local usernames under HKLM\SAM\SAM\Domains\Account\Users \Names
  - Passwords are under Users key above Names

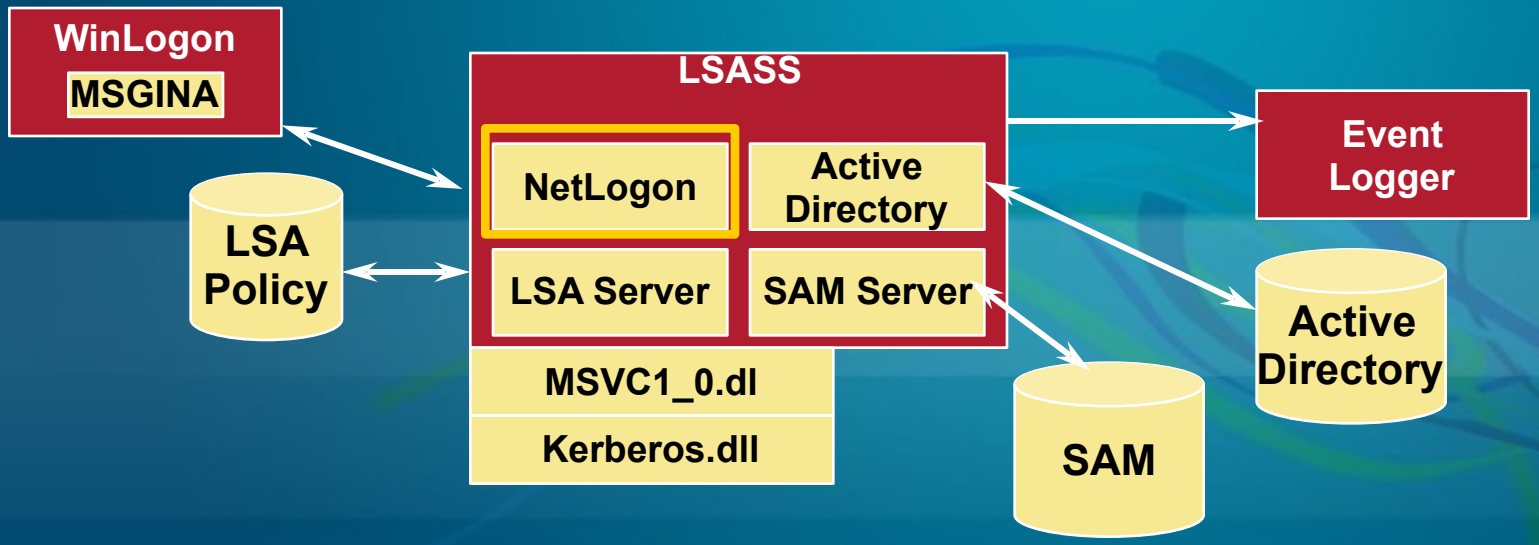

Administrator: Command Prompt
C:\Users\Administrator>psexec -s -i -d c:\Windows\regedit.exe
PsExec v2.2 - Execute processes remotely
Copyright (C) 2001-2016 Mark Russinovich
Sysinternals - www.sysinternals.com
c:\Windows\regedit.exe started on WIN-JWBPPZSXEF
Registry Editor
File Edit View Favorites Help
Computer
HKEY_CLASSES_ROOT
HKEY_CURRENT_USER
HKEY_LOCAL_MACHINE
BCD00000000
COMPONENTS
HARDWARE
SAM
SAM
Domains
Account
Aliases
Groups
Users
000001F4
000001F5
000003E8
Names
Administrator
Guest
student
Builtin
RXACT

LSASS Components

- Active Directory
  - A directory service that contains a database that stores information about objects in a domain
  - A *domain* is a collection of computers and their associated security groups that are managed as a single entity
  - The Active Directory server, implemented as a service, \Windows\System32\Ntdsa.dll, that runs in the Lsass process
- Authentication packages
  - DLLs that run in the context of the Lsass process and that implement Windows authentication policy:
    - LanMan: \Windows\System32\Msv1_0.dll
    - Kerberos: \Windows\System32\Kerberos.dll
    - Negotiate: uses LanMan or Kerberos, depending on which is most appropriate

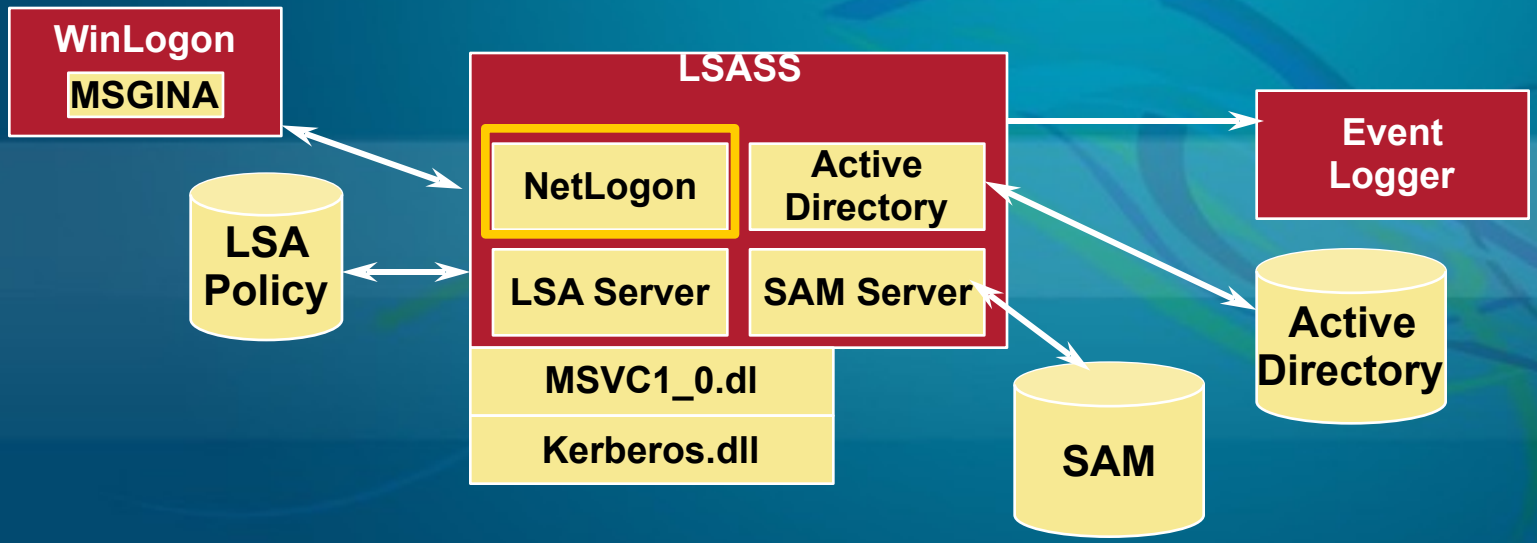
LSASS Components

- Net Logon service (Netlogon)
  - A Windows service (\Windows\System32\Netlogon.dll) that runs inside Lsass and responds to Microsoft LAN Manager 2 Windows NT (pre-Windows 2000) network logon requests
  - Authentication is handled as local logons are, by sending them to Lsass for verification
  - Netlogon also has a locator service built into it for locating domain controllers



Winlogon

- Logon process (Winlogon)
  - A user-mode process running \Windows\System32\Winlogon.exe that is responsible for responding to the SAS and for managing interactive logon sessions
- Graphical Identification and Authentication (GINA)
  - A user-mode DLL that runs in the Winlogon process and that Winlogon uses to obtain a user's name and password or smart card PIN
    - Default is \Windows\System32\Msgina.dll



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What Makes Logon Secure?

- Before anyone logs on, the visible desktop is Winlogon's
- Winlogon registers CTRL+ALT+DEL, the Secure Attention Sequence (SAS), as a standard hotkey sequence
- SAS takes you to the Winlogon desktop
- No application can deregister it because only the thread that registers a hotkey can deregister it
- When Windows' keyboard input processing code sees SAS it disables keyboard hooks so that no one can intercept it

Logon

- After getting security identification (account name, password), the GINA sends it to the Local Security Authority Sub System (LSASS)
- LSASS calls an authentication package to verify the logon
  - If the logon is local or to a legacy domain, MSV1_0 is the authenticator. User name and password are encrypted and compared against the Security Accounts Manager (SAM) database
    - Cached domain logons are also handled by MSV1_0
  - If the logon is to a AD domain the authenticator is Kerberos, which communicates with the AD service on a domain controller
- If there is a match, the SIDs of the corresponding user account and its groups are retrieved
- Finally, LSASS retrieves account privileges from the Security database or from AD

Logon

- LSASS creates a token for your logon session and Winlogon attaches it to the first process of your session
  - Tokens are created with the NtCreateToken API
  - Every process gets a copy of its parent's token
- SIDs and privileges cannot be added to a token
- A logon session is active as long as there is at least one token associated with the session
- Lab
  - Run "LogonSessions –p" (from Sysinternals) to view the active logon sessions on your system

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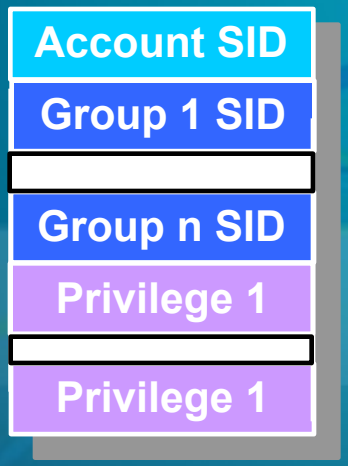
The Access Validation Algorithm

- Access validation is a security equation that takes three inputs:
  - Desired Access
  - Process Token
    - Or Thread's token if the thread is "impersonating"
  - The object's Security Descriptor, which contains a Discretionary Access Control List (DACL)
- The output is access allowed or access denied

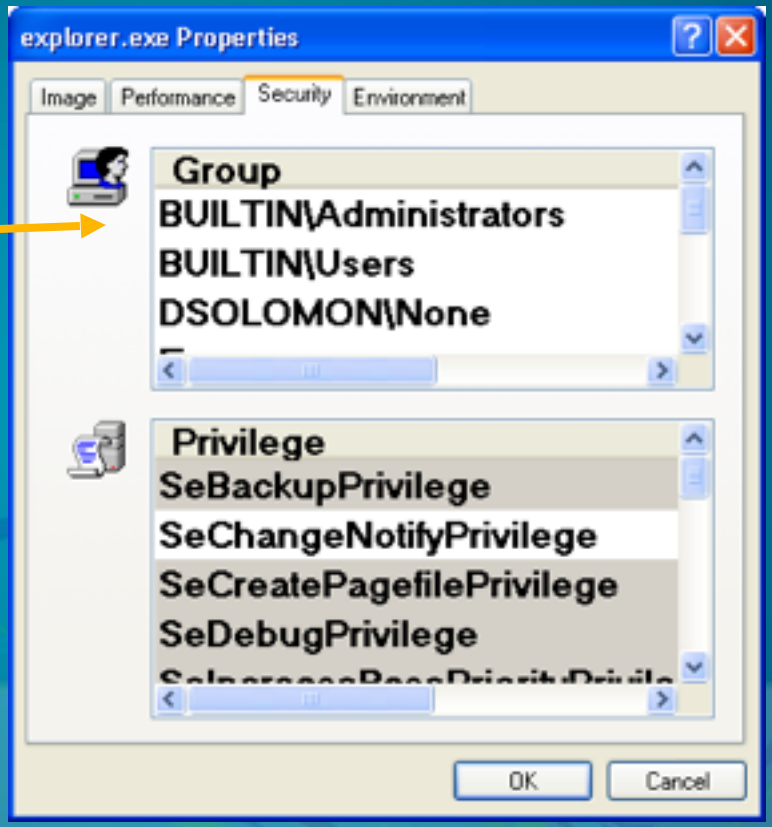
Tokens

- The main components of a token are:
  - SID of the user
  - SIDs of groups the user account belongs to
  - Privileges assigned to the user (described in next section)



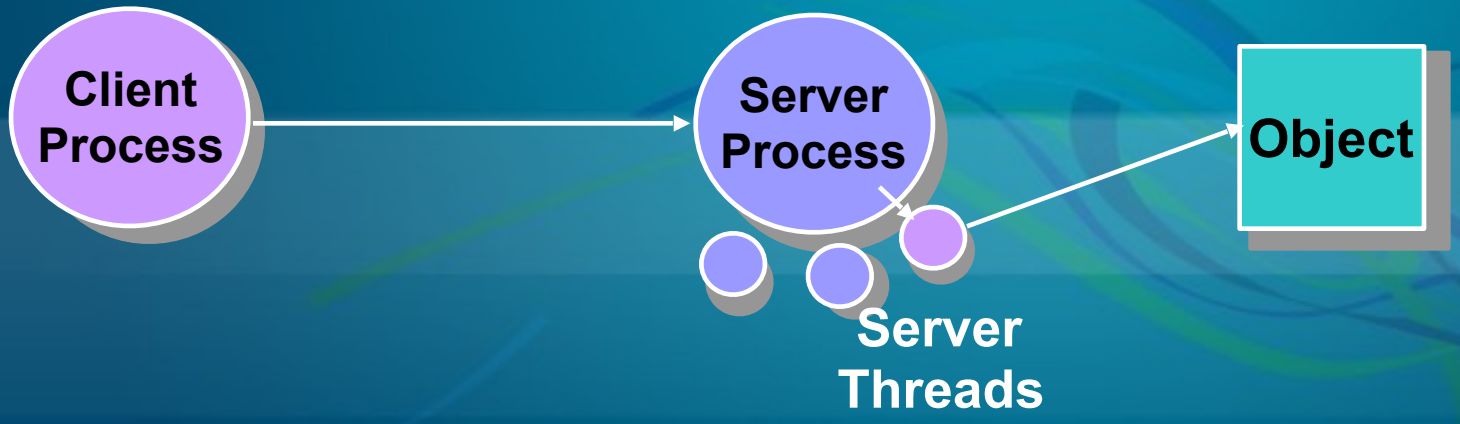
Labs: Viewing Access Tokens

- Process Explorer: double click on a process and go to Security tab
  - Examine groups list
- Use RUNAS to create a CMD process running under another account (e.g. your domain account)
  - Examine groups list
- Viewing tokens with the Kernel Debugger
  - Run !process 0 0 to find a process
  - Run !process <PID> 1 to dump the process
  - Get the token address and type !token –n <token address>
  - Type dt _token <token address> to see all fields defined in a token

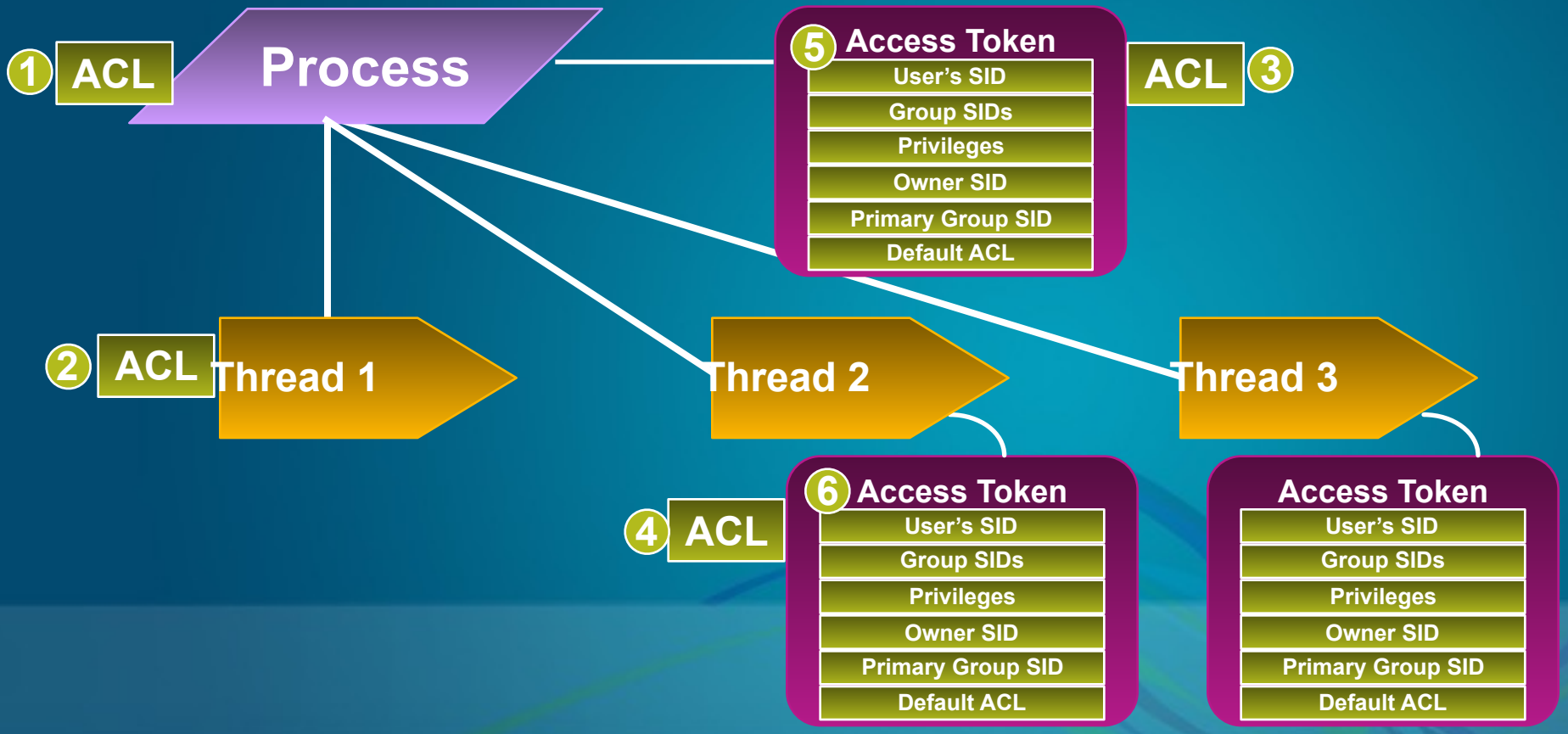


Impersonation

- Lets an application adopt the security profile another user
  - Used by server applications
  - Impersonation is implemented at the thread level
    - The process token is the “primary token” and is always accessible
    - Each thread can be impersonating a different client
- Can impersonate with a number of client/server networking APIs – named pipes, RPC, DCOM



Process And Thread Security Structures



- Thread tokens (where present) completely supersede process token (basis for "security impersonation")

SIDs

- Windows uses Security Identifers (SIDs) to identify security principles:
  - Users, Groups of users, Computers, Domains
- SIDs consist of:
  - A revision level e.g. 1
  - An identifier-authority value e.g. 5 (SECURITY_NT_AUTHORITY)
  - One or more subauthority values
- Who assigns SIDs?
  - Setup assigns a computer a SID
  - Dcpromo assigns a domain a SID
  - Users and groups on the local machine are assigned SIDs that are rooted with the computer SID, with a Relative Identifier (RID) at the end
    - RIDs start at 1000 (built-in account RIDs are pre-defined)
- Some local users and groups have pre-defined SIDs (eg. World = S-1-1-0)

Demo: SIDs

- Example SIDs

Domain SID: S-1-5-21-34125455-5125555-1251255

First account: S-1-5-21-34125455-5125555-1251255-1000

Admin account: S-1-5-21-34125455-5125555-1251255-500

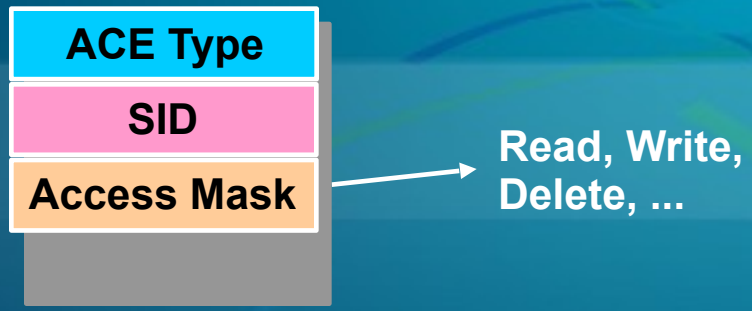
System account: S-1-5-18

```
C:\Users\Administrator>whoami /all

USER INFORMATION
----------------

User Name                      SID
============================== ============================================
win-jwbppzsxefv\administrator S-1-5-21-1367486129-1636748403-2738611465-500
```

DACLs

- DACLs consist of zero or more Access Control Entries
  - A security descriptor with no DACL allows all access
  - A security descriptor with an empty (0-entry) DACL denies everybody all access

- An ACE is either “allow” or “deny”



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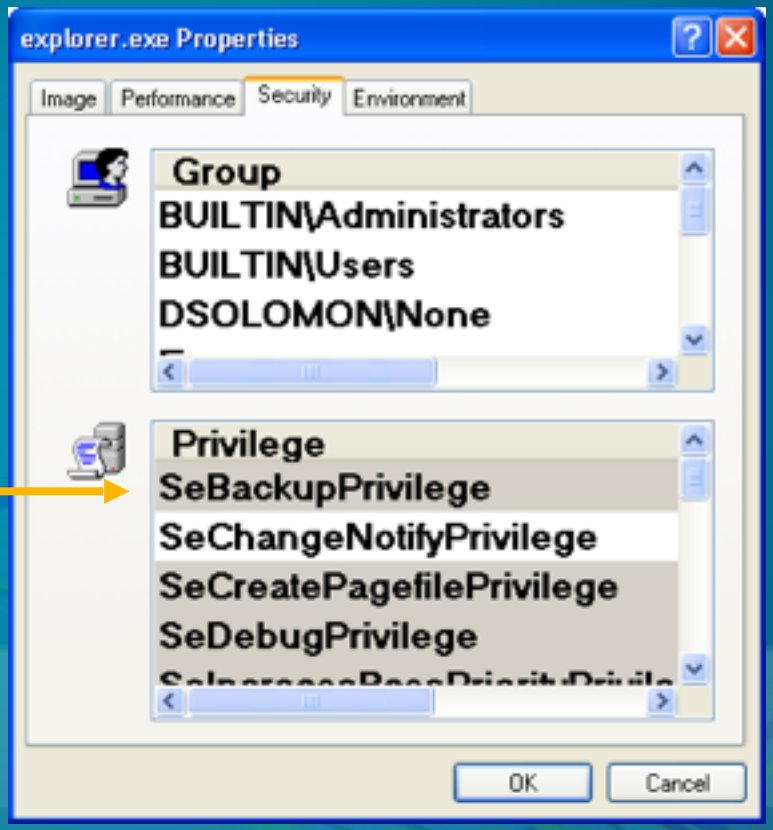
Privileges

- Specify which system actions a process (or thread) can perform
- Privileges are associated with groups and user accounts
  - There are sets of pre-defined privileges associated with built-in groups (e.g. System, Administrators)
- Examples include:
  - Backup/Restore
  - Shutdown
  - Debug
  - Take ownership
- Privileges are disabled by default and must be programmatically turned on with a system call



Demo: Privileges

- Run Secpol.msc and examine full list
  - Click on Local Policies->User Rights assignment
- Process Explorer: double click on a process, go to security tab, and examine privileges list
- Watch changes to privilege list:
  1. Run Process Explorer – put in paused mode
  2. Open Control Panel applet to change system time
  3. Go back to Process Explorer & press F5
  4. Examine privilege list in new process that was created
  5. Notice in privilege list that system time privilege is enabled

Powerful Privileges

- There are several privileges that gives an account that has them full control of a computer:
  - Debug: can open any process, including System processes to
    - Inject code
    - Modify code
    - Read sensitive data
  - Take Ownership: can access any object on the system
    - Replace system files
    - Change security
  - Restore: can replace any file
  - Load Driver
    - Drivers bypass all security
  - Create Token
    - Can spoof any user (locally)
    - Requires use of undocumented NT API
  - Trusted Computer Base (Act as Part of Operating System)
    - Can create a new logon session with arbitrary SIDs in the token

Demo: Powerful Privileges

- View the use of the backup privilege:
  - Make a directory
  - Create a file in the directory
  - Use the security editor to remove inherited security and give Everyone full access to the file
  - Remove all access to the directory (do not propagate)
  - Start a command-prompt and do a “dir” of the directory
  - Run \Sysint\Solomon\PView and enable the Backup privilege for the command prompt
  - Do another “dir” and note the different behavior
- View the use of the Bypass-Traverse Checking privilege (internally called “Change Notify”)
  - From the same command prompt run notepad to open the file (give the full path) in the inaccessible directory
  - Extra credit: disable Bypass-Traverse Checking so that you get access denied trying to open the file (hint: requires use of secpol.msc and then RUNAS)